

Memory, Data, & Addressing II

Roadmap

C:

```
car *c = malloc(sizeof(car));  
c->miles = 100;  
c->gals = 17;  
float mpg = get_mpg(c);  
free(c);
```

Java:

```
Car c = new Car();  
c.setMiles(100);  
c.setGals(17);  
float mpg =  
    c.getMPG();
```

Memory & data

Arrays & structs

Integers & floats

RISC V assembly

Procedures & stacks

Executables

Memory & caches

Processor Pipeline

Performance

Parallelism

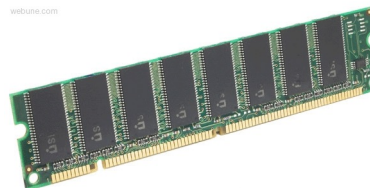
Assembly
language:

```
get_mpg(car*):  
    lw    a5,0(a0)  
    lw    a4,4(a0)  
    divw  a5,a5,a4  
    fcvf.s.w    fa0,a5  
    ret
```

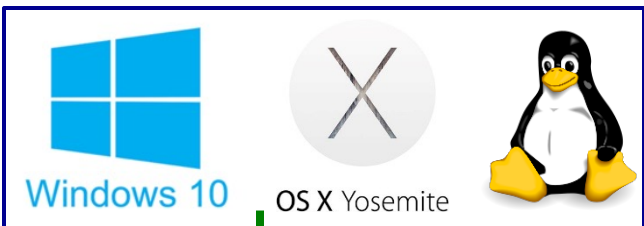
Machine
code:

```
0111010000011000  
100011010000010000000010  
1000100111000010  
110000011111101000011111
```

Computer
system:



OS:



Review

- 1) If the word size of a machine is 64-bits, which of the following is usually true? (pick all that apply)
 - a) 64 bits is the size of a pointer
 - b) 64 bits is the size of an integer
 - c) 64 bits is the width of a register
- 2) (True/False) By looking at the bits stored in memory, I can tell if a particular 4-bytes is being used to represent an integer, floating point number, or instruction.
- 3) If the size of a pointer on a machine is 6 bits, the address space is how many bytes?

Memory, Data, and Addressing

- ❖ Representing information as bits and bytes
- ❖ Organizing and addressing data in memory
- ❖ **Manipulating data in memory using C**
- ❖ Boolean algebra and bit-level manipulations

Addresses and Pointers in C

❖ $\&$ = “address of” operator

❖ $*$ = “value at address” or “dereference” operator

$*$ is also used with variable declarations

```
int* ptr;
```

Declares a variable, `ptr`, that is a pointer to (i.e. holds the address of) an `int` in memory

```
int x = 5;
```

```
int y = 2;
```

Declares two variables, `x` and `y`, that hold `ints`, and *initializes* them to 5 and 2, respectively

```
ptr = &x;
```

Sets `ptr` to the address of `x` (“`ptr` points to `x`”)

```
y = 1 + *ptr;
```

“Dereference `ptr`”

Sets `y` to “1 plus the value stored at the address held by `ptr`. Because `ptr` points to `x`, this is equivalent to `y=1+x`;

What is `*(&y)` ?

Assignment in C

- ❖ A variable is represented by a location
- ❖ Declaration $\neq$ initialization (initially holds “garbage”)
- ❖ `int x, y;`
 - `x` is at address `0x04`, `y` is at `0x18`

	0x00	0x01	0x02	0x03	
0x00	A7	00	32	00	
0x04	00	01	29	F3	x
0x08	EE	EE	EE	EE	
0x0C	FA	CE	CA	FE	
0x10	26	00	00	00	
0x14	00	00	10	00	
0x18	01	00	00	00	y
0x1C	FF	00	F4	96	
0x20	DE	AD	BE	EF	
0x24	00	00	00	00	

Assignment in C

32-bit example
(pointers are 32-bits wide)

little-endian

- ❖ A variable is represented by a location
- ❖ Declaration $\neq$ initialization (initially holds “garbage”)
- ❖ `int x, y;`
 - `x` is at address `0x04`, `y` is at `0x18`

	0x00	0x01	0x02	0x03	
0x00					
0x04	00	01	29	F3	x
0x08					
0x0C					
0x10					
0x14					
0x18	01	00	00	00	y
0x1C					
0x20					
0x24					

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"

* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

❖ `int x, y;`

❖ `x = 0;`

	0x00	0x01	0x02	0x03	
0x00					x
0x04	00	00	00	00	
0x08					
0x0C					
0x10					y
0x14					
0x18	01	00	00	00	
0x1C					
0x20					
0x24					

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"

* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

❖ **int** x, y;

❖ x = 0;

❖ y = 0x3CD02700;

little endian!

	0x00	0x01	0x02	0x03	
0x00					
0x04	00	00	00	00	x
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	y
0x1C					
0x20					
0x24					

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"

* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

❖ **int** x, y;

❖ x = 0;

❖ y = 0x3CD02700;

❖ x = y + 3;

- Get value at y, add 3, store in x

	0x00	0x01	0x02	0x03	
0x00					x
0x04	03	27	D0	3C	
0x08					
0x0C					
0x10					y
0x14					
0x18	00	27	D0	3C	
0x1C					
0x20					
0x24					

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"

* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

❖ **int** x, y;

❖ x = 0;

❖ y = 0x3CD02700;

❖ x = y + 3;

- Get value at y, add 3, store in x

❖ **int*** z;

- z is at address 0x20

	0x00	0x01	0x02	0x03	
0x00					
0x04	03	27	D0	3C	x
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	y
0x1C					
0x20	DE	AD	BE	EF	z
0x24					

initial value is whatever bits were already there! ("garbage")

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"

* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

❖ **int** x, y;

❖ x = 0;

❖ y = 0x3CD02700;

❖ x = y + 3;

- Get value at y, add 3, store in x

❖ **int*** z = &y + 3;

- Get address of y, "add 3", store in z

	0x00	0x01	0x02	0x03	
0x00					
0x04	03	27	D0	3C	x
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	y
0x1C					
0x20	24	00	00	00	z
0x24					

Pointer arithmetic



Pointer Arithmetic

- ❖ Pointer arithmetic is scaled by the size of target type
 - In this example, `sizeof(int) = 4`
- ❖ `int* z = &y + 3;`
 - Get address of `y`, add `3*sizeof(int)`, store in `z`
 - $\&y = 0x18 = 1*16^1 + 8*16^0 = 24$
 - $24 + 3*(4) = 36 = 2*16^1 + 4*16^0 = 0x24$
- ❖ **Pointer arithmetic can be dangerous!**
 - Can easily lead to bad memory accesses
 - Be careful with data types and *casting*

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"

* = "dereference"

- ❖ `int x, y;`
- ❖ `x = 0;`
- ❖ `y = 0x3CD02700;`
- ❖ `x = y + 3;`
 - Get value at `y`, add 3, store in `x`
- ❖ `int* z = &y + 3;`
 - Get address of `y`, add **12**, store in `z`
- ❖ `*z = y;`
 - What does this do?

	0x00	0x01	0x02	0x03	
0x00					
0x04	03	27	D0	3C	x
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	y
0x1C					
0x20	24	00	00	00	z
0x24					

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"

* = "dereference"

❖ `int x, y;`

❖ `x = 0;`

❖ `y = 0x3CD02700;`

❖ `x = y + 3;`

■ Get value at `y`, add 3, store in `x`

❖ `int* z = &y + 3;`

■ Get address of `y`, add **12**, store in `z`

The target of a pointer
is also a location

❖ `*z = y;`

■ Get value of `y`, put in address
stored in `z`

	0x00	0x01	0x02	0x03	
0x00					
0x04	03	27	D0	3C	x
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	y
0x1C					
0x20	24	00	00	00	z
0x24	00	27	D0	3C	

Arrays in C

Declaration: **int** **a**[6];

element type

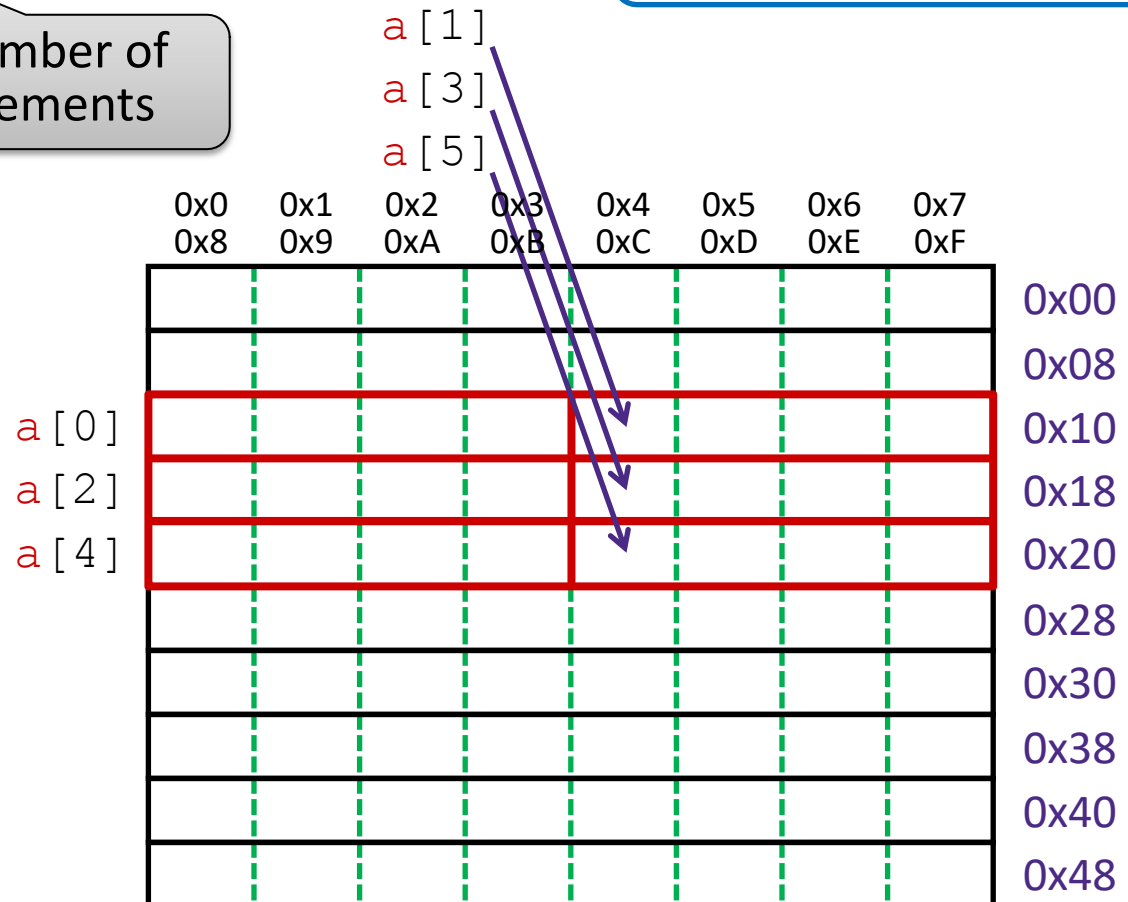
name

number of
elements

Arrays are adjacent locations in memory storing the same type of data object

a (array name) returns the array's address

64-bit example
(pointers are 64-bits wide)



Arrays Basics

- **Pitfall:** An array in C does not know its own length, and its bounds are not checked!
 - We can accidentally access off the end of an array
 - We must pass the array **and its size** to any procedure that is going to manipulate it
- Mistakes with array bounds cause *segmentation faults* and *bus errors*
 - Be careful! These are VERY difficult to find (You'll learn how to debug these in lab)

Arrays in C

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes

	0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF	
									0x00
									0x08
<code>a[0]</code>	5F	01	00	00					0x10
<code>a[2]</code>									0x18
<code>a[4]</code>					5F	01	00	00	0x20
									0x28
									0x30
									0x38
									0x40
									0x48

Arrays in C

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes

	0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF	
									0x00
					AD	0B	00	00	0x08
<code>a[0]</code>	5F	01	00	00					0x10
<code>a[2]</code>									0x18
<code>a[4]</code>					5F	01	00	00	0x20
	AD	0B	00	00					0x28
									0x30
									0x38
									0x40
									0x48

Arrays in C

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes

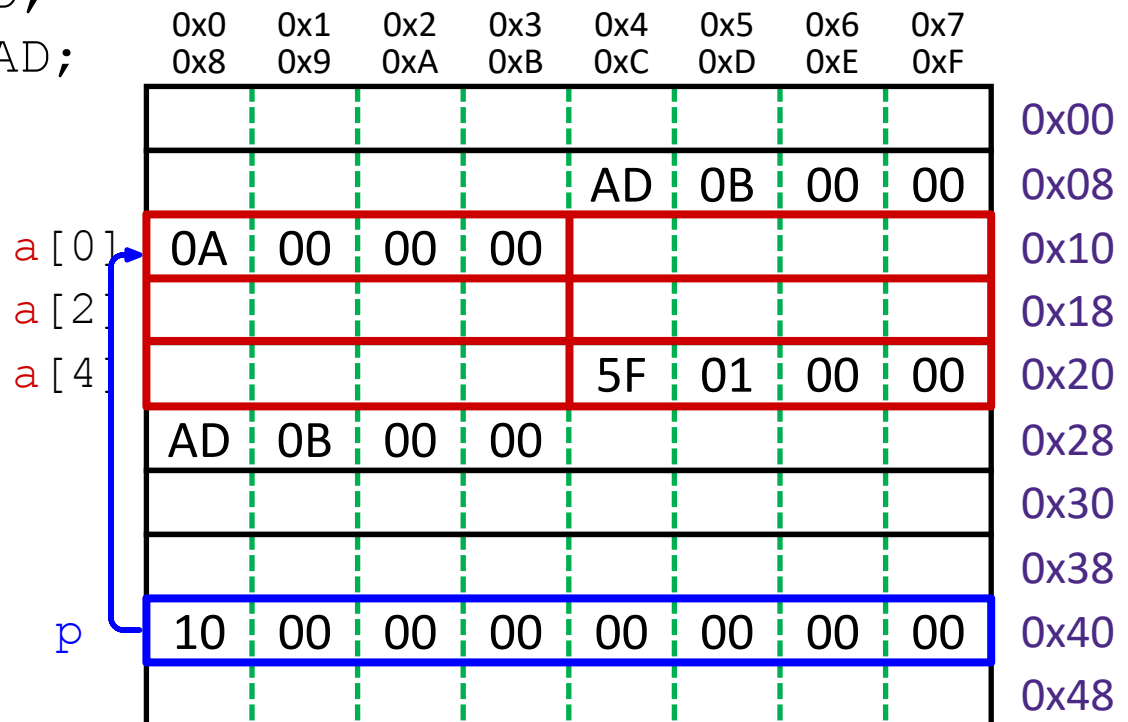
Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`

equivalent $\left\{ \begin{array}{l} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{array} \right.$



Arrays in C

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`

equivalent $\left\{ \begin{array}{l} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{array} \right.$

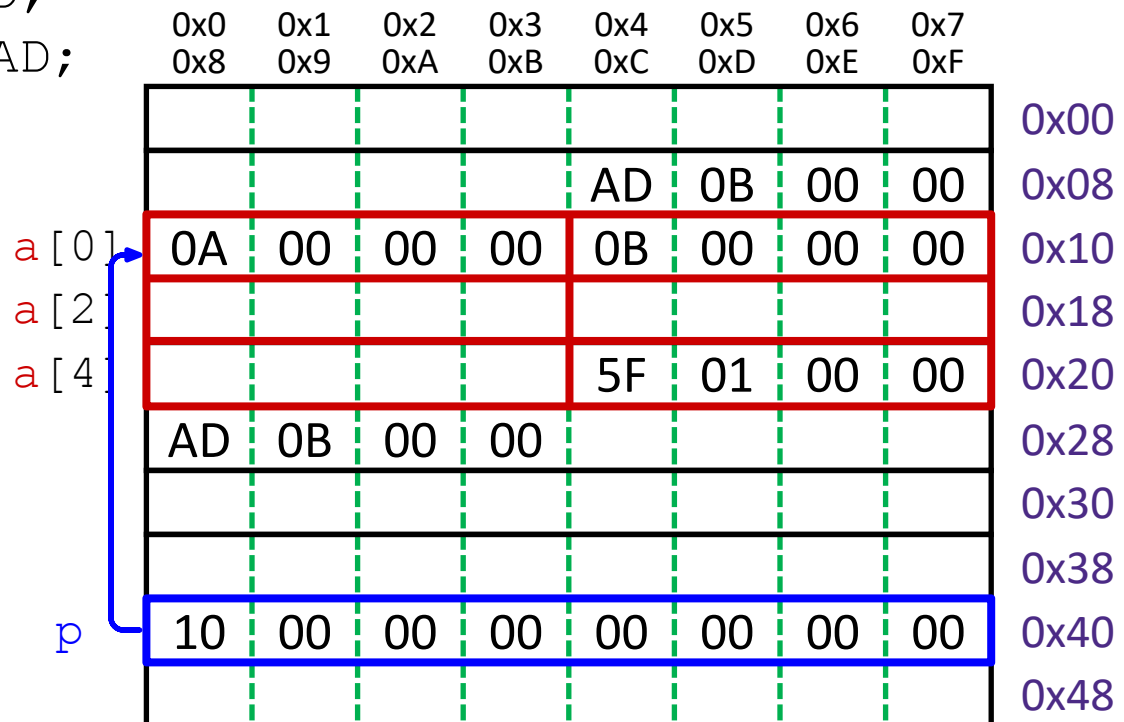
array indexing = address arithmetic
 (both scaled by the size of the type)

equivalent $\left\{ \begin{array}{l} p[1] = 0xB; \\ *(p+1) = 0xB; \\ p = p + 2; \end{array} \right.$

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes



Arrays in C

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`

equivalent $\left\{ \begin{array}{l} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{array} \right.$

array indexing = address arithmetic
 (both scaled by the size of the type)

equivalent $\left\{ \begin{array}{l} p[1] = 0xB; \\ *(p+1) = 0xB; \\ p = p + 2; \end{array} \right.$

`*p = a[1] + 1;`

	0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF	
					AD	0B	00	00	0x00
									0x08
<code>a[0]</code>	0A	00	00	00	0B	00	00	00	0x10
<code>a[2]</code>	0C	00	00	00					0x18
<code>a[4]</code>					5F	01	00	00	0x20
	AD	0B	00	00					0x28
									0x30
									0x38
<code>p</code>	18	00	00	00	00	00	00	00	0x40
									0x48

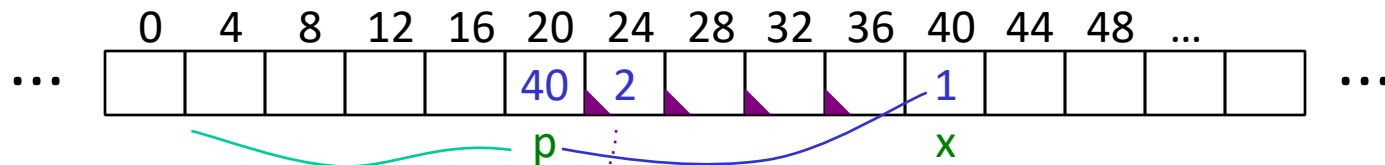
Arrays Stored Differently Than Pointers

```

→ void foo() {
→   int *p, a[4], x;
→   p = &x;

→   *p = 1; // or p[0]
→   printf("*p:%u, p:%u, &p:%u\n", *p, p, &p);
→   *a = 2; // or a[0]
→   printf("*a:%u, a:%u, &a:%u\n", *a, a, &a);
→ }

```



***p:1, p:40, &p:20**

***a:2, a:24, &a:24**

K&R: "An array name is not a variable"

Representing strings

- ❖ C-style string stored as an array of bytes (**char***)
 - Elements are one-byte **ASCII codes** for each character
 - No “String” keyword, unlike Java

32	space	48	0	64	@	80	P	96	`	112	p
33	!	49	1	65	A	81	Q	97	a	113	q
34	"	50	2	66	B	82	R	98	b	114	r
35	#	51	3	67	C	83	S	99	c	115	s
36	\$	52	4	68	D	84	T	100	d	116	t
37	%	53	5	69	E	85	U	101	e	117	u
38	&	54	6	70	F	86	V	102	f	118	v
39	'	55	7	71	G	87	W	103	g	119	w
40	(	56	8	72	H	88	X	104	h	120	x
41	)	57	9	73	I	89	Y	105	i	121	y
42	*	58	:	74	J	90	Z	106	j	122	z
43	+	59	;	75	K	91	[	107	k	123	{
44	,	60	<	76	L	92	\	108	l	124	
45	-	61	=	77	M	93	]	109	m	125	}
46	.	62	>	78	N	94	^	110	n	126	~
47	/	63	?	79	O	95	_	111	o	127	del

ASCII: American Standard Code for Information Interchange

Null-Terminated Strings

- ❖ **Example:** "Donald Trump" stored as a 13-byte array

<i>Decimal:</i>	68	111	110	97	108	100	32	84	114	117	109	112	0
<i>Hex:</i>	0x44	0x6F	0x6E	0x61	0x6C	0x64	0x20	0x54	0x72	0x75	0x6D	0x70	0x00
<i>Text:</i>	D	o	n	a	l	d		T	r	u	m	p	\0

- ❖ Last character followed by a 0 byte (' \0 ')
(a.k.a. "**null terminator**")
 - Must take into account when allocating space in memory
 - Note that ' 0 ' $\neq$ ' \0 ' (*i.e.* character 0 has non-zero value)
- ❖ How do we compute the length of a string?
 - Traverse array until null terminator encountered

C (char = 1 byte)

Endianness and Strings

```
char s[6] = "12345";
```

String literal

0x31 = 49 decimal = ASCII '1'

IA32, x86-64

(little-endian)

SPARC

(big-endian)

0x00	31	↔	31	0x00	'1'
0x01	32	↔	32	0x01	'2'
0x02	33	↔	33	0x02	'3'
0x03	34	↔	34	0x03	'4'
0x04	35	↔	35	0x04	'5'
0x05	00	↔	00	0x05	'\0'

- ❖ Byte ordering (endianness) is not an issue for 1-byte values
 - The whole array does not constitute a single value
 - Individual elements are values; chars are single bytes

Examining Data Representations

❖ Code to print byte representation of data

- Any data type can be treated as a *byte array* by **casting** it to `char`
- C has **unchecked** casts **!! DANGER !!**

```
void show_bytes(char* start, int len) {  
    int i;  
    for (i = 0; i < len; i++)  
        printf("%p\t0x%.2x\n", start+i, *(start+i));  
    printf("\n");  
}
```

printf directives:

<code>%p</code>	Print pointer
<code>\t</code>	Tab
<code>%x</code>	Print value as hex
<code>\n</code>	New line

Examining Data Representations

❖ Code to print byte representation of data

- Any data type can be treated as a *byte array* by **casting** it to `char`
- C has **unchecked** casts **!! DANGER !!**

```
void show_bytes(char* start, int len) {  
    int i;  
    for (i = 0; i < len; i++)  
        printf("%p\t0x%.2x\n", start+i, *(start+i));  
    printf("\n");  
}
```

```
void show_int(int x) {  
    show_bytes( (char *) &x, sizeof(int));  
}
```

show_bytes Execution Example

```
int x = 12345; // 0x00003039
printf("int x = %d;\n", x);
show_int(x); // show_bytes((char *) &x, sizeof(int));
```

❖ Result (Linux x86-64):

- **Note:** The addresses will change on each run (try it!), but fall in same general range

```
int x = 12345;
0x7fffb7f71dbc      0x39
0x7fffb7f71dbd      0x30
0x7fffb7f71dbe      0x00
0x7fffb7f71dbf      0x00
```

Summary

- ❖ Assignment in C results in value being put in memory location
- ❖ Pointer is a C representation of a data address
 - `&` = “address of” operator
 - `*` = “value at address” or “dereference” operator
- ❖ Pointer arithmetic scales by size of target type
 - Convenient when accessing array-like structures in memory
 - Be careful when using – particularly when *casting* variables
- ❖ Arrays are adjacent locations in memory storing the same type of data object
 - Strings are null-terminated arrays of characters (ASCII)

Assignment in C - Handout

32-bit example
(pointers are 32-bits wide)

& = "address of"

* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value*
 - Store RHS value at LHS location

❖ **int** x, y;

❖ x = 0;

❖ y = 0x3CD02700;

❖ x = y + 3;

❖ **int*** z = &y + 3;

❖ *z = y;

0x00

0x04

0x08

0x0C

0x10

0x14

0x18

0x1C

0x20

0x24

0x00 0x01 0x02 0x03

x

y

z

Arrays in C - Handout

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`

equivalent $\left\{ \begin{array}{l} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{array} \right.$

`a[0]`
`a[2]`
`a[4]`

array indexing = address arithmetic
 (both scaled by the size of the type)

equivalent $\left\{ \begin{array}{l} p[1] = 0xB; \\ *(p+1) = 0xB; \\ p = p + 2; \end{array} \right.$

`*p = a[1] + 1;`

	0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF
0x00								
0x08								
0x10								
0x18								
0x20								
0x28								
0x30								
0x38								
0x40								
0x48								