

Integers I



<http://xkcd.com/257/>

Acknowledgments: These slides have been modified by Arrvinth Shriraman, Justin Tsia

Roadmap

C:

```
car *c = malloc(sizeof(car));
c->miles = 100;
c->gals = 17;
float mpg = get_mpg(c);
free(c);
```

Java:

```
Car c = new Car();
c.setMiles(100);
c.setGals(17);
float mpg =
    c.getMPG();
```

Memory & data
Arrays & structs
Integers & floats
RISC V assembly
Procedures & stacks
Executables
Memory & caches
Processor Pipeline
Performance
Parallelism

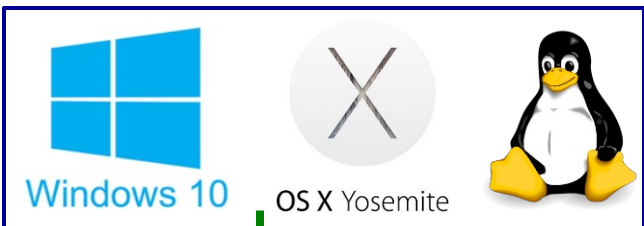
Assembly
language:

```
get_mpg(car*):
    lw    a5,0(a0)
    lw    a4,4(a0)
    divw  a5,a5,a4
    fcvf.s.w    fa0,a5
    ret
```

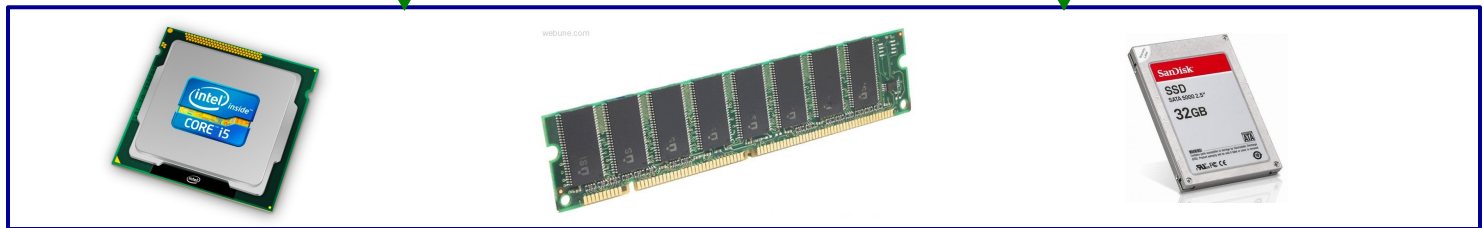
Machine
code:

```
0111010000011000
100011010000010000000010
1000100111000010
110000011111101000011111
```

OS:

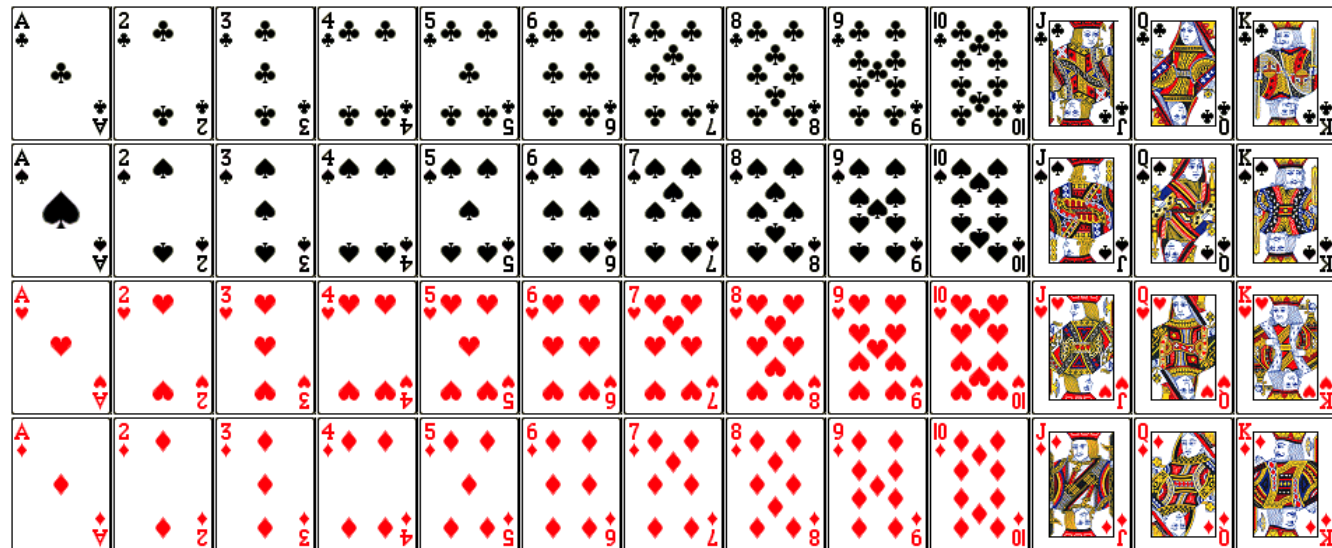


Computer
system:



But before we get to integers....

- ❖ Encode a standard deck of playing cards
 - How do we encode suits, face cards?
- ❖ 52 cards in 4 suits
 - Which is the higher value card?
 - Are they the same suit?



Boolean Algebra

- ❖ Developed by George Boole in 19th Century
 - Algebraic representation of logic (True $\rightarrow$ 1, False $\rightarrow$ 0)
 - AND: $A \& B = 1$ when both A is 1 and B is 1
 - OR: $A | B = 1$ when either A is 1 or B is 1
 - XOR: $A \wedge B = 1$ when either A is 1 or B is 1, but not both
 - NOT: $\sim A = 1$ when A is 0 and vice-versa
 - DeMorgan's Law:

$$\sim (A | B) = \sim A \& \sim B$$

$$\sim (A \& B) = \sim A | \sim B$$

AND			OR			XOR			NOT		
$\&$	0	1		0	1		0	1	$\sim$		
0	0	0	0	0	1	0	0	1	0	1	
1	0	1	1	1	1	1	1	0	1	0	

General Boolean Algebras

❖ Operate on bit vectors

- Operations applied bitwise
- All of the properties of Boolean algebra apply

01101001	01101001	01101001	
& 01010101	01010101	^ 01010101	~ 01010101

❖ Examples of useful operations:

$$x \wedge x = 0$$

01010101
^ 01010101
00000000

$$x | 1 = 1, \quad x | 0 = x$$

01010101
11110000
11110101

Bit-Level Operations in C

❖ $\&$ (AND), $|$ (OR), $\wedge$ (XOR), $\sim$ (NOT)

- View arguments as bit vectors, apply operations bitwise
- Apply to any “integral” data type
 - long, int, short, char, unsigned

❖ Examples with char a, b, c;

- `a = (char) 0x41;` `// 0x41->0b 0100 0001`
`b = ~a;` `//                      0b                      ->0x`
- `a = (char) 0x69;` `// 0x69->0b 0110 1001`
`b = (char) 0x55;` `// 0x55->0b 0101 0101`
`c = a & b;` `//                      0b                      ->0x`
- `a = (char) 0x41;` `// 0x41->0b 0100 0001`
`b = a;` `//                      0b 0100 0001`
`c = a ^ b;` `//                      0b                      ->0x`

Contrast: Logic Operations

❖ Logical operators in C: `&&` (AND), `||` (OR), `!` (NOT)

- 0 is False, anything nonzero is True
- Always return 0 or 1
- **Early termination** (a.k.a. short-circuit evaluation) of `&&`, `||`
 - `int x = (42 == 6) || boom();`
 - `int y = (42 == 6) && boom();`

❖ Examples (char data type)

- | | |
|----------------------------------|--|
| ■ <code>!0x41 -> 0x00</code> | ■ <code>0xCC && 0x33 -> 0x01</code> |
| ■ <code>!0x00 -> 0x01</code> | ■ <code>0x00 0x33 -> 0x01</code> |
| ■ <code>!!0x41 -> 0x01</code> | |

Two possible representations

1) 1 bit per card (52): bit corresponding to card set to 1



- “One-hot” encoding (similar to set notation)
- Drawbacks:
 - Hard to compare values and suits
 - Large number of bits required

2) 1 bit per suit (4), 1 bit per number (13): 2 bits set



- Pair of one-hot encoded values
- Easier to compare suits and values, but still lots of bits used

Two better representations

3) Binary encoding of all 52 cards – only 6 bits needed

- $2^6 = 64 \geq 52$



low-order 6 bits of a byte

- Fits in one byte (smaller than one-hot encodings)
- How can we make value and suit comparisons easier?

4) Separate binary encodings of suit (2 bits) and value (4 bits)



suit value

- Also fits in one byte, and easy to do comparisons

K	Q	J	...	3	2	A
1101	1100	1011	...	0011	0010	0001

	00
	01
	10
	11

Compare Card Suits

mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector v .
Here we turns all *but* the bits of interest in v to 0.

```
char hand[5];           // represents a 5-card hand
char card1, card2;      // two cards to compare
card1 = hand[0];
card2 = hand[1];
...
if ( isSameSuit(card1, card2) ) { ... }
```

```
#define SUIT_MASK 0x30
```

```
int isSameSuit(char card1, char card2) {
    return (card1 & SUIT_MASK) == (card2 & SUIT_MASK);
}
```

returns int

SUIT_MASK = 0x30 =

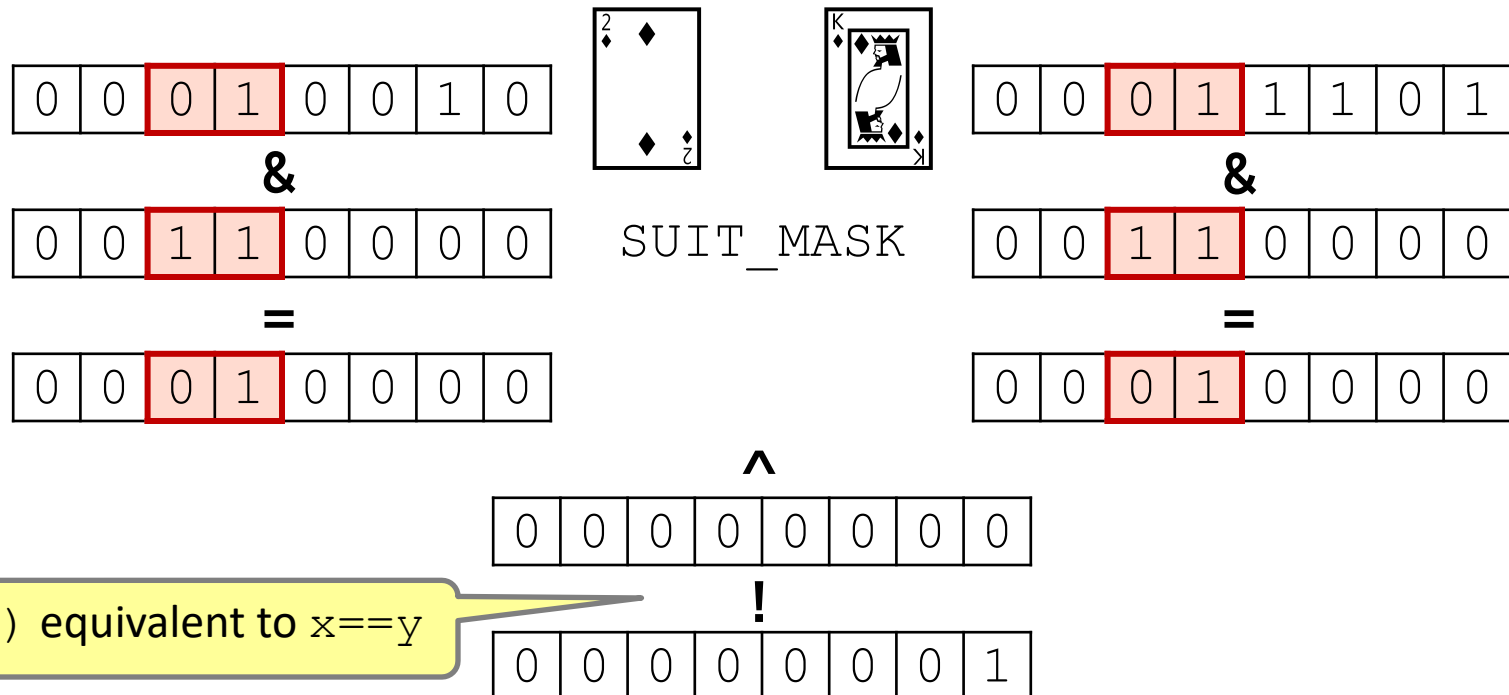
0	0	1	1	0	0	0	0
suit				value			

Compare Card Suits

mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector v .
Here we turns all *but* the bits of interest in v to 0.

```
#define SUIT_MASK 0x30
```

```
int isSameSuit(char card1, char card2) {  
    return (!((card1 & SUIT_MASK) ^ (card2 & SUIT_MASK)));  
    // return (card1 & SUIT_MASK) == (card2 & SUIT_MASK);  
}
```



Compare Card Values

mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector v .

```
char hand[5];           // represents a 5-card hand
char card1, card2;      // two cards to compare
card1 = hand[0];
card2 = hand[1];

...

if ( greaterValue(card1, card2) ) { ... }
```

```
#define VALUE_MASK 0x0F

int greaterValue(char card1, char card2) {
    return ((unsigned char)(card1 & VALUE_MASK) >
            (unsigned char)(card2 & VALUE_MASK));
}
```

VALUE_MASK = 0x0F =

0	0	0	0	1	1	1	1
---	---	---	---	---	---	---	---

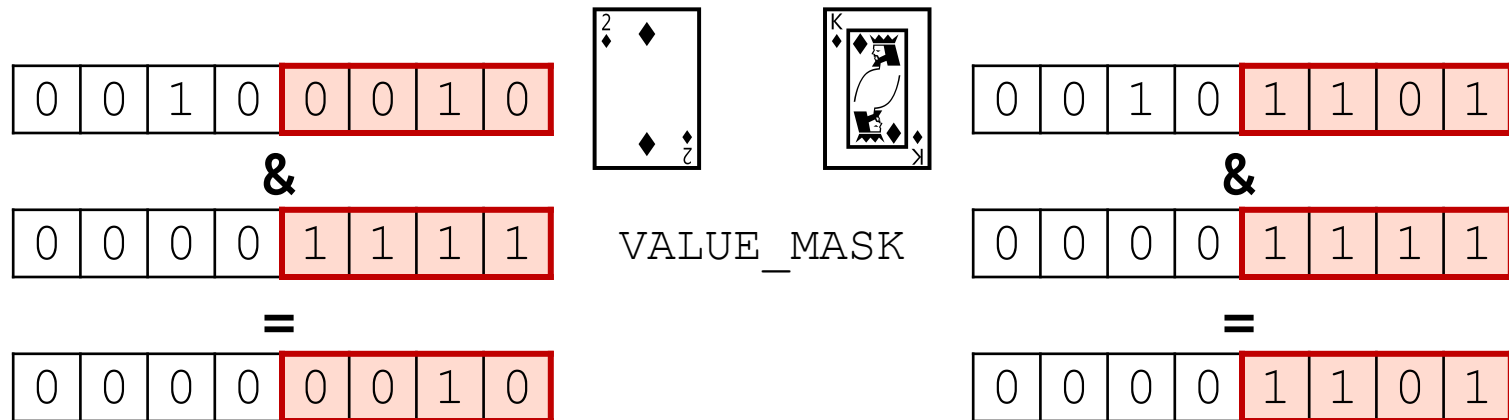
suit
value

Compare Card Values

mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector v .

```
#define VALUE_MASK 0x0F

int greaterValue(char card1, char card2) {
    return ((unsigned int)(card1 & VALUE_MASK) >
            (unsigned int)(card2 & VALUE_MASK));
}
```



$2_{10} > 13_{10}$

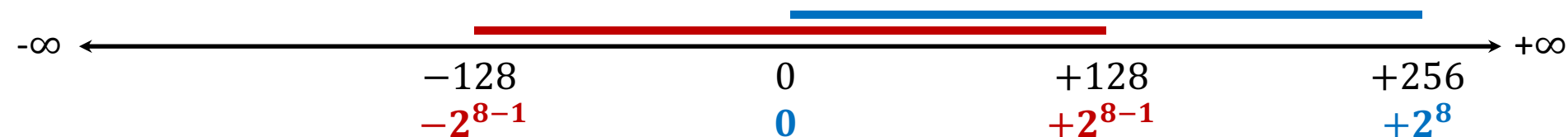
0 (false)

Integers

- ❖ **Binary representation of integers**
 - **Unsigned and signed**
 - Casting in C
- ❖ Consequences of finite width representation
 - Overflow, sign extension
- ❖ Shifting and arithmetic operations

Encoding Integers

- ❖ The hardware (and C) supports two flavors of integers
 - *unsigned* – only the non-negatives
 - *signed* – both negatives and non-negatives
- ❖ Cannot represent all integers with w bits
 - Only 2^w distinct bit patterns
 - Unsigned values: $0 \dots 2^w - 1$
 - Signed values: $-2^{(w-1)} \dots 0 \dots 2^{(w-1)} - 1$
- ❖ **Example:** 8-bit integers (*e.g.* `char`)



Unsigned Integers

- ❖ Unsigned values follow the standard base 2 system
 - $b_7b_6b_5b_4b_3b_2b_1b_0 = b_72^7 + b_62^6 + \dots + b_12^1 + b_02^0$
- ❖ Add and subtract using the normal “carry” and “borrow” rules, just in binary

63
+ 8
71

00111111
+00001000
01000111

- ❖ Useful formula: $N \text{ ones in a row} = 2^N - 1$
- ❖ How would you make *signed* integers?

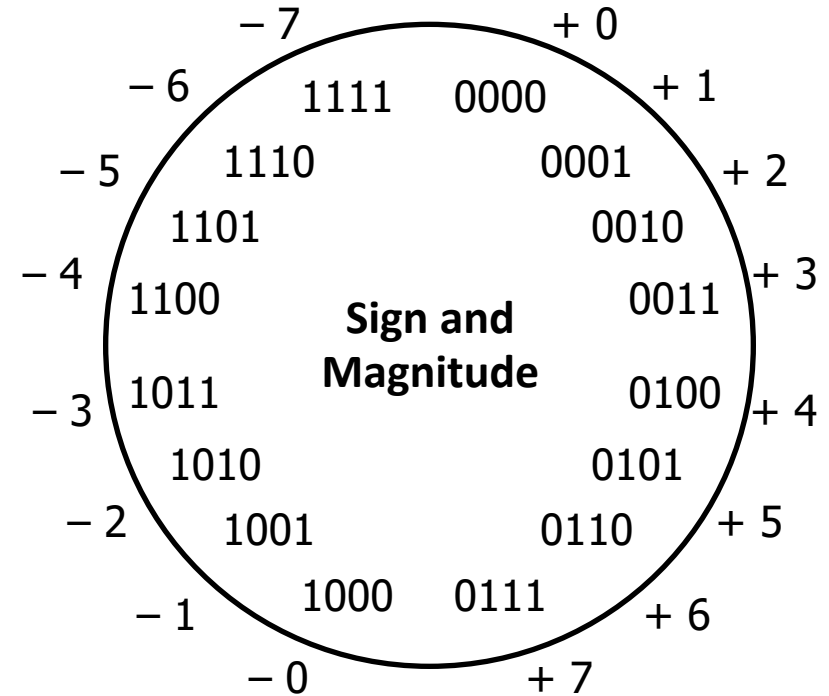
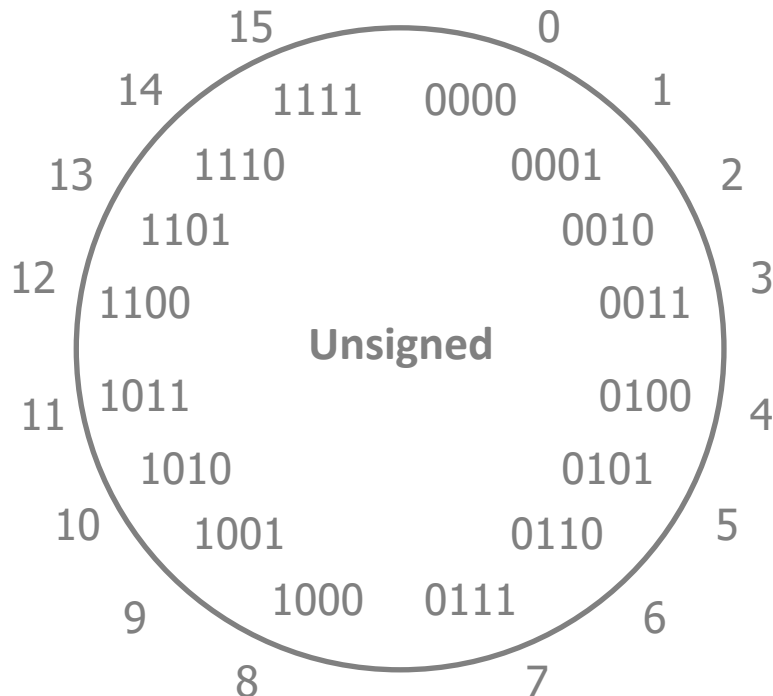
Sign and Magnitude

Most Significant Bit

- ❖ Designate the high-order bit (MSB) as the “sign bit”
 - `sign=0`: positive numbers; `sign=1`: negative numbers
- ❖ Benefits:
 - Using MSB as sign bit matches positive numbers with unsigned
 - All zeros encoding is still = 0
- ❖ Examples (8 bits):
 - $0x00 = 00000000_2$ is non-negative, because the sign bit is 0
 - $0x7F = 01111111_2$ is non-negative ($+127_{10}$)
 - $0x85 = 10000101_2$ is negative (-5_{10})
 - $0x80 = 10000000_2$ is negative... zero???

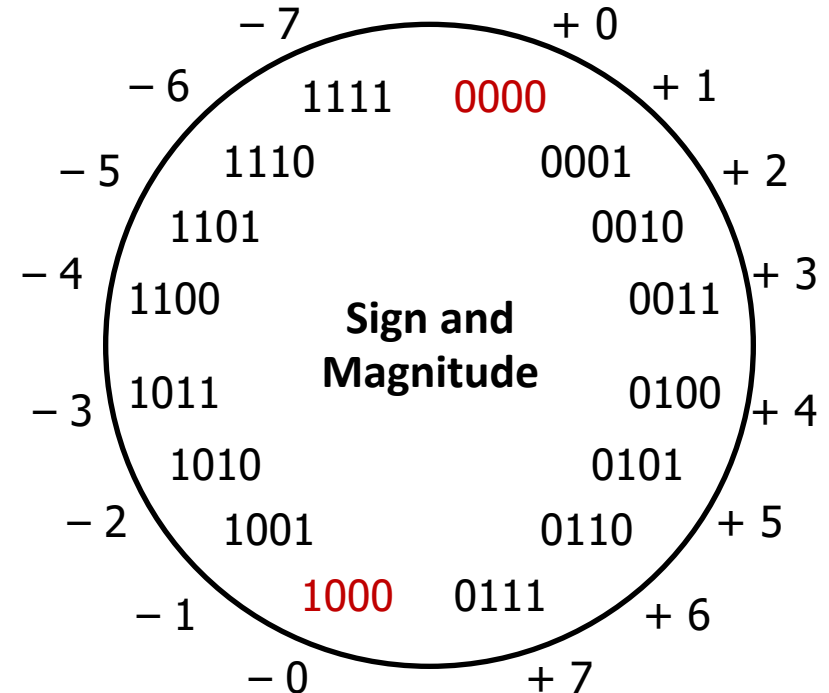
Sign and Magnitude

- ❖ MSB is the sign bit, rest of the bits are magnitude
- ❖ Drawbacks?



Sign and Magnitude

- ❖ MSB is the sign bit, rest of the bits are magnitude
- ❖ Drawbacks:
 - **Two representations of 0** (bad for checking equality)



Sign and Magnitude

❖ MSB is the sign bit, rest of the bits are magnitude

❖ Drawbacks:

■ Two representations of 0 (bad for checking equality)

■ **Arithmetic is cumbersome**

• Example: $4 - 3 \neq 4 + (-3)$

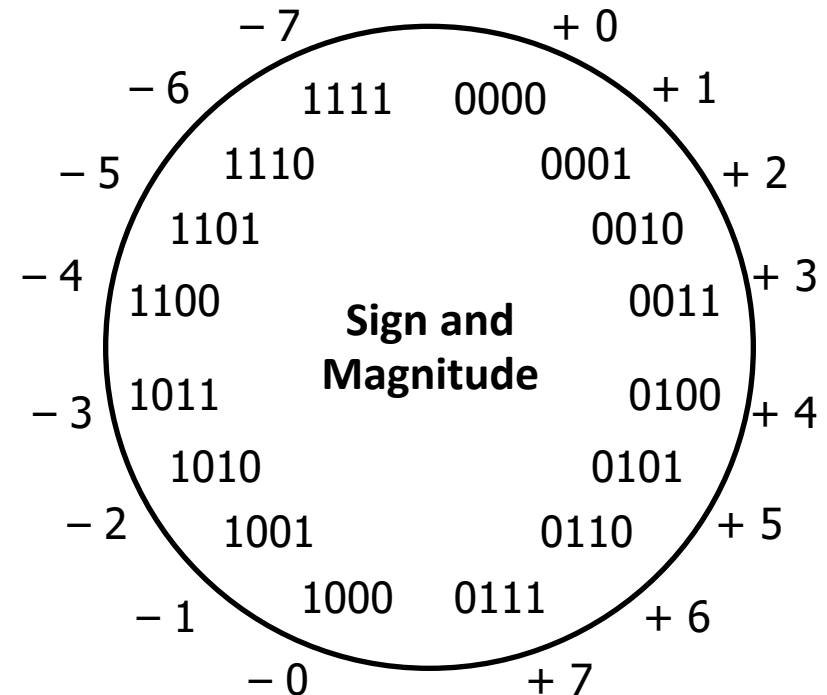
4	0100
- 3	- 0011
1	0001



4	0100
+ -3	+ 1011
-7	1111



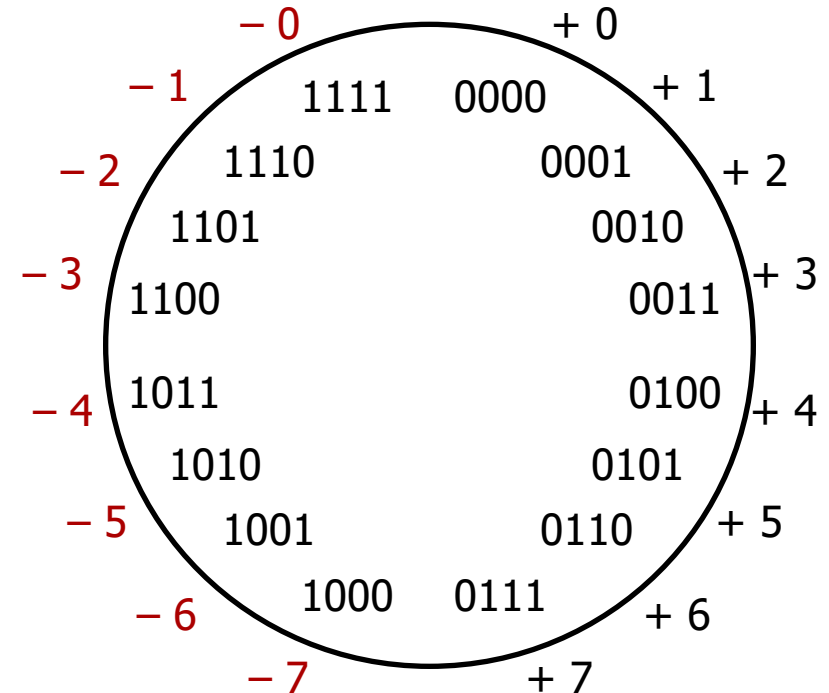
• Negatives “increment” in wrong direction!



Two's Complement

❖ Let's fix these problems:

1) “Flip” negative encodings so incrementing works



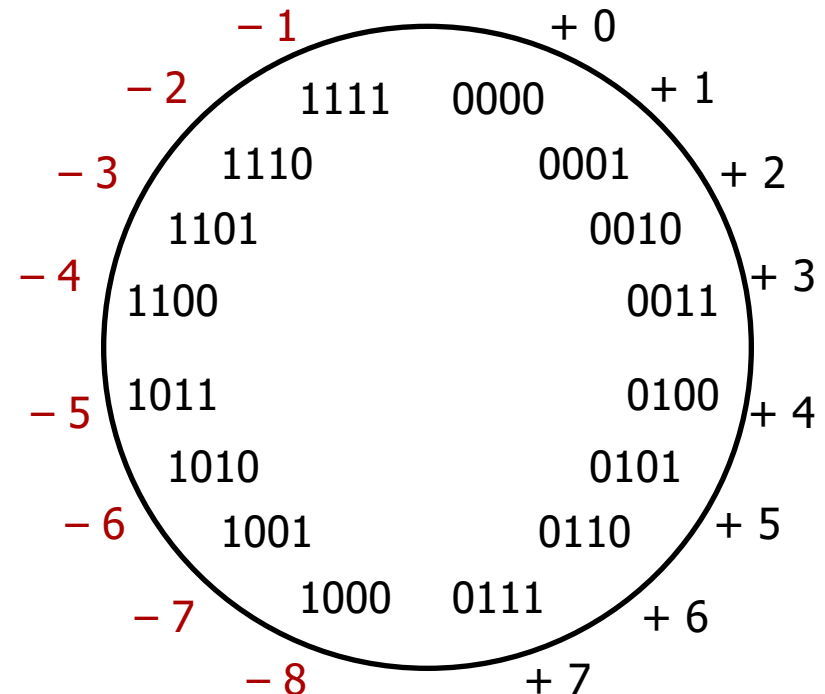
Two's Complement

❖ Let's fix these problems:

- 1) “Flip” negative encodings so incrementing works
- 2) “Shift” negative numbers to eliminate -0

❖ MSB *still* indicates sign!

- This is why we represent one more negative than positive number (-2^{N-1} to $2^{N-1} - 1$)



Two's Complement Negatives

- ❖ Accomplished with one neat mathematical trick!

$b_{(w-1)}$ has weight $-2^{(w-1)}$, other bits have usual weights $+2^i$



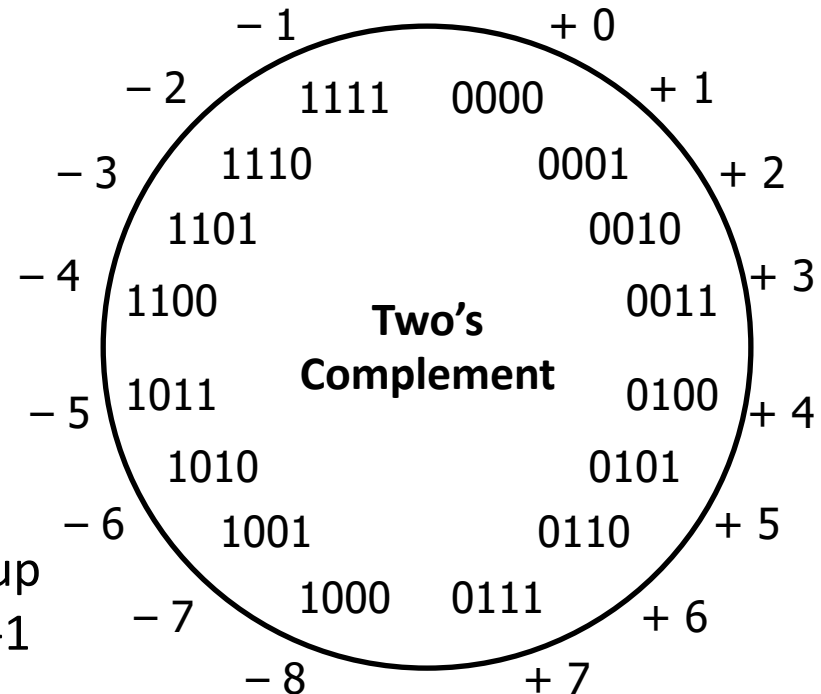
4-bit Examples:

- 1010_2 unsigned:
 $1*2^3 + 0*2^2 + 1*2^1 + 0*2^0 = 10$
- 1010_2 two's complement:
 $-1*2^3 + 0*2^2 + 1*2^1 + 0*2^0 = -6$

-1 represented as:

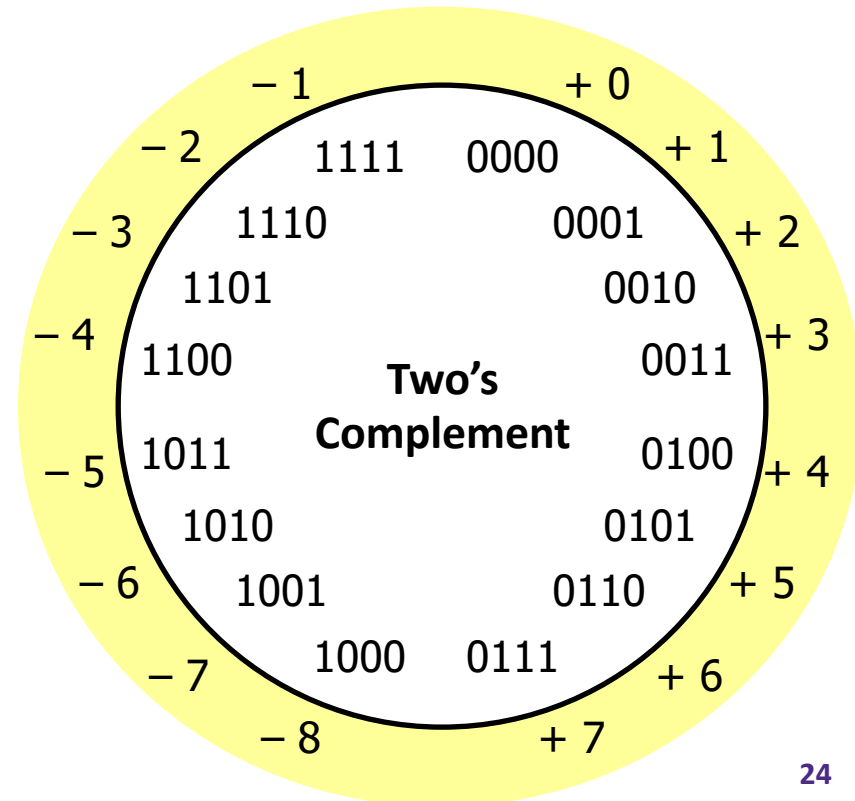
$$1111_2 = -2^3 + (2^3 - 1)$$

- MSB makes it super negative, add up all the other bits to get back up to -1



Why Two's Complement is So Great

- ❖ Roughly same number of (+) and (−) numbers
 - ❖ Positive number encodings match unsigned
 - ❖ Simple arithmetic ($x + -y = x - y$)
 - ❖ Single zero
 - ❖ All zeros encoding = 0
 - ❖ Simple negation procedure:
 - Get negative representation of any integer by taking bitwise complement and then adding one!
- (~x + 1 == -x)



Summary

- ❖ Bit-level operators allow for fine-grained manipulations of data
 - Bitwise AND (`&`), OR (`|`), and NOT (`~`) **different than logical** AND (`&&`), OR (`||`), and NOT (`!`)
 - Especially useful with bit masks
- ❖ Choice of *encoding scheme* is important
 - Tradeoffs based on size requirements and desired operations
- ❖ Integers represented using unsigned and two's complement representations
 - Limited by fixed bit width
 - We'll examine arithmetic operations next lecture