

CMPT 225 D100, Fall 2025

Final Exam December 5, 2025

Name _____

SFU ID: |_|_|_|_|_|_|_|_|_|_|

Problem 1	
Problem 2	
Problem 3	
Problem 4	
TOTAL	

Instructions:

1. The duration of the exam is 180 minutes.
2. Write your full name and SFU ID ****clearly****.
3. This is a closed book exam, no calculators, cell phones, or any other material.
4. The exam consists of four (4) problems, each worth 25 points
5. Write your answers in the provided space.
6. There is an extra page at the end of the exam. You may use it if needed.
7. You may write helper functions if needed.
8. Explain all your answers.
9. Really, explain all your answers.

Good luck!

Problem 1 (Java syntax, Big-O notation) [25 points]

A. (5 points) Prove that the running time of the function `fun_A(n)` is $O(n \log(n))$.

```
int fun_A (int n) {  
    int sum = 0;  
    for (int i=0; i<n; i++)  
        for (int j=1; j<i; j=j*2)  
            sum += 1;  
    return sum;  
}
```

B. (5 points) Prove that the running time of the function `fun_B(n)` is $O(n)$.

```
int fun_B (int n) {  
    int sum = 0;  
    for (int i=1; i<n; i*=2)  
        for (int j=1; j<i; j++)  
            sum += j;  
    return sum;  
}
```

Consider the function `fun_C(a,b)`.

```
static void printAB(int a, int b) {  
    for (int i = a; i <= b; i++)  
        System.out.print(i + " ");  
    System.out.println();  
}  
static void fun_C(int a, int b) {  
    if (a < b) {  
        printAB(a, b);  
        int q = (b-a) / 2;  
        fun_C(a, a+q);  
        fun_C(a+q+1, b);  
    }  
}
```

C. (7 points) What will be the output of `fun_C(2,6)`? List all recursive calls of the program

D. (8 points) What is the running time of `fun_C(0,n)` as a function of n ?
Use big-O notation to express your answer.

Problem 2 (Graphs) [25 points]

A. (5 points) Draw the graph defined by this adjacency matrix.

	A	B	C	D	E	F	G
A	0	1	0	1	0	0	0
B	1	0	0	1	1	0	0
C	0	0	0	0	1	0	1
D	1	1	0	0	0	0	0
E	0	1	1	0	0	1	1
F	0	0	0	0	1	0	0
G	0	0	1	0	1	0	0

B. (20 points) A map is given as a $n \times n$ 2d array of ints. The cells of the map are connected *horizontally/vertically* (not diagonally).

Zeros correspond to *water* and positive numbers correspond to *land*.

Each island consists of connected cells of land (positive numbers).

For example, suppose grid = {

```
{5,0,0,0,0,0},
{4,0,2,5,0,0},
{0,0,1,0,0,0},
{0,3,2,4,0,0},
{0,6,0,0,2,0},
{0,0,0,1,8,0}};
```

This map has 3 islands, colored with **blue**, **red**, and **green**. Note that the green and red islands contain cells that touch on a diagonal, and we do not consider it as connected.

Write a function that gets a $n \times n$ map, and coordinates (x,y) , where $0 \leq x,y < n$.

- If `map[x][y] == 0`, the function returns 0.
- Otherwise, it returns the sum of all numbers on the island containing (x,y) .

In the example above,

- The sum on the blue island is $5+4=9$.
- The sum on the red island is $2+5+1+3+2+4+6=23$.
- The sum on the green island is $2+1+8=11$.

For full marks the running time of the function should be $O(\text{size of the island})$

Explain your answer before writing code.

```
public int sumIsland(int[][] map, int n, int x, int y) {
```

Problem 3 (Binary Trees) [25 points]

In this problem use the following definition of Binary Tree. You may assume the classes have the standard getters/setters.

```
public class BTreeNode<T> {  
    private T data;  
    private BTreeNode<T> leftChild;  
    private BTreeNode<T> rightChild;  
    private BTreeNode<T> parent;  
}  
  
public class BinaryTree<T> {  
    private BTreeNode<T> root;  
}
```

- A. (10 points) Write the method `equals()` for the class Binary Tree. The method returns true if `this` has the exact same tree structure as the `other`, with the same data in the corresponding nodes. Compare if two nodes contain the same data using `==`.
Explain the idea of the algorithm and its running time.

```
public boolean equals(BinaryTree<T> other) {
```

- B. (10 points) Write the method `nextInorder()`. The method gets a node in a binary tree and returns the next node in the inorder traversal of the tree. The running time must be at most $O(\text{height of the tree})$

Explain the idea of the algorithm and its running time.

```
public BTNode<T> nextInorder(BTNode<T> node) {
```

- C. (5 points) Draw a **Binary Search Tree** whose preOrder traversal is [5,1,2,3,4,8,7,6,9].

Problem 4 (Heaps) [25 points]

- A. (10 points) A min-heap is typically represented using an array. Write the following functions for min-heap. Explain your answer (e.g., by drawing a correspondence between the tree and the array)

1. *// returns the index in the array containing the root*
`public int` getRoot() {

}
2. *// returns the index in the array containing the left child of the node i*
`public int` getLeftChild(`int` i) {

}
3. *// returns the index in the array containing the right child of the node i*
`public int` getRightChild(`int` i) {

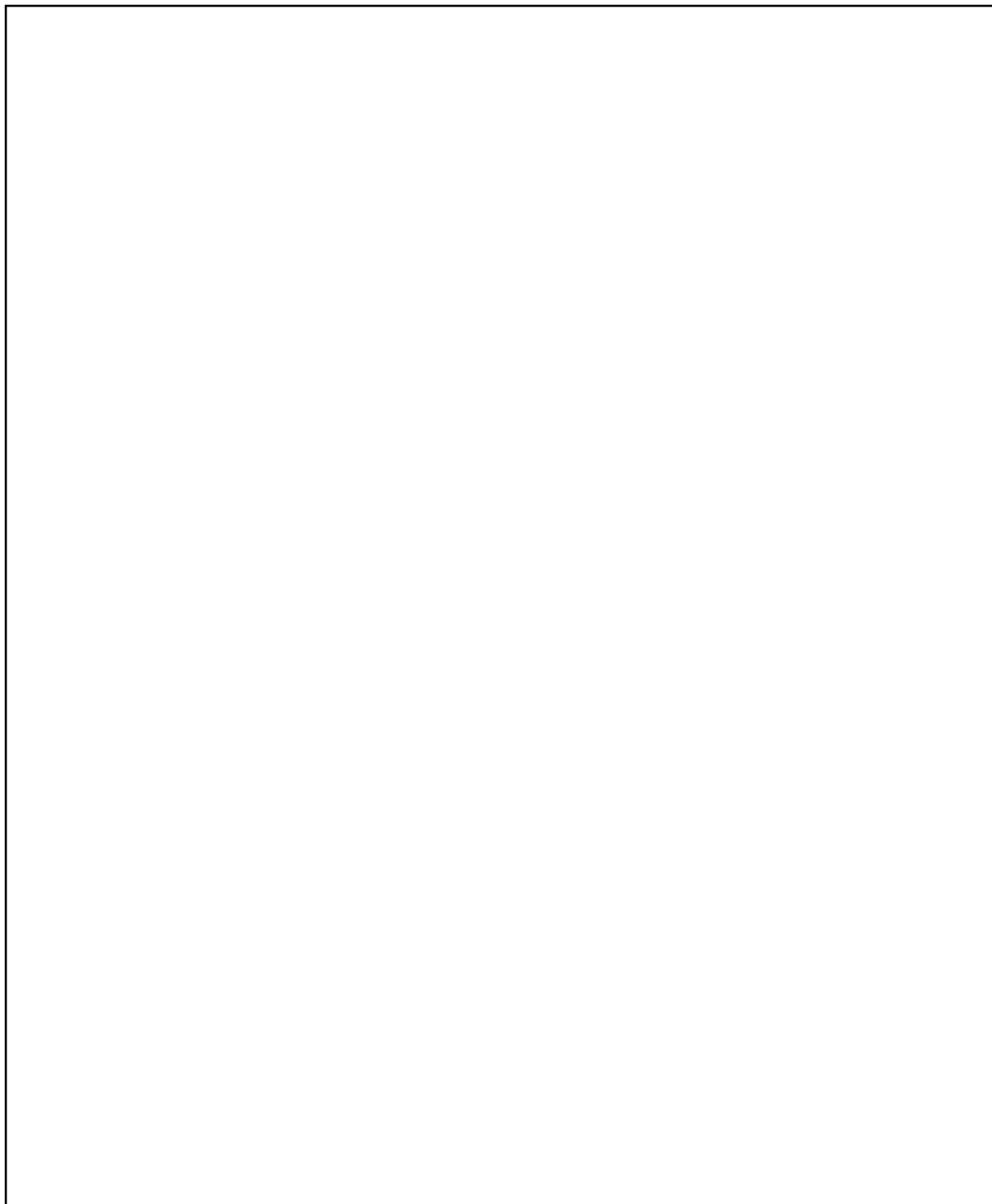
}
4. *// returns the index in the array containing the parent of the node i*
`public int` getParent(`int` i) {

}

- B. (8 points) Consider the array of integers $A = [1, 10, 15, 20, 50, 30, 35, 80, 2, 3, 4, 5, 6]$.
Is there a *min-heap* whose preOrder traversal is A ?
If yes, draw the min-heap. If not, explain why.

- C. (7 points) Apply **buildMinHeap** on the array $[9, 8, 7, 6, 5, 4, 3, 2, 1]$.
Show some intermediate steps by drawing the corresponding tree and the array.
Recall buildHeap works as follows:
- ```
public void buildHeap(int[] ar) {
 for (int i=ar.length-1; i>=0; i--)
 heapify(ar,i);
}
```

**Extra page**





**Master Theorem**

Let  $T(n) = a T(n/b) + f(n)$ ,  $T(1) = O(1)$

Define  $c = \log_b(a)$

- If  $f(n) = O(n^d)$  for  $d < c$ , then  $T(n) = \Theta(n^c)$
- If  $f(n) = \Omega(n^d)$  for  $d > c$ , then  $T(n) = \Theta(n^d)$
- If  $f(n) = \Theta(n^c)$ , then  $T(n) = \Theta(n^c \log(n))$