CMPT 125, Fall 2021

Final Exam - Solution December 17, 2021

Name	 	 	 	 	
SFU ID:				 	l

Instructions:

- 1. Duration of the exam is 180 minutes.
- 2. Write your name and SFU ID **clearly**.
- 3. This is a closed book exam, no calculators, cell phones, or any other material.
- 4. The exam consists of four (4) problems. Each problem is worth 25 points.
- 5. Write your answers in the provided space.
- 6. There is an extra page at the end of the exam. You may use it if needed.
- 7. Explain all your answers.
- 8. Really, explain all your answers.

Good luck!

Problem 1 [25 points]

a) The function below gets as input an array of ints of length n...

```
void what(int* ar, int n) {
    if (n>1) {
        what(ar, n-1);
        char tmp = ar[n-2];
        ar[n-2] = ar[n-1];
        ar[n-1] = tmp;
    }
}
```

[6 points] What is the time complexity of what() as a function of n?. Use Big-O notation to express your answer. Explain your solution.

Answer: time complexity on an array of length n can be written as T(n) = O(1) + T(n-1). This behaves like T(n) = O(n)

[6 points] What is the effect of what() when applied on the array [1,2,3,4,5]? What is the functionality of what()? Explain your answer.

Answer: let's try to do do it backwards:

- On input [1] the array stays the same
- On input [1,2] the array becomes [2,1].
- On input [1,2,3] we make a recursive call on [1,2] and change it to [2,1], and then swap 3 and 1. So the result is [2,3,1].
- On input [1,2,3,4] we make a recursive call on [1,2,3] and change it to [2,3,1], and then swap 4 and 1. So the result is [3,2,4,1].
- On input [1,2,3,4,5] we make a recursive call on [1,2,3,4] and change it to [2,3,4,1], and then swap 4 and 1. So the result is [2,3,4,5,1].

In general, the function rotates the given array by one to the left. (Same logic as above)

```
b) [7 points] What will be the output of the C++ program below? Explain your answer.
#include <iostream>
using namespace std;
class Test {
  private:
    int m x;
  public:
    Test() {
      m x=0;
    Test(int x) {
     m \times = \times;
    Test(Test& other) {
     this->m x = other.m x+1;
    int getX() { return m x; }
    void setX(int x) { m x=x; }
};
int main() {
  Test t; calls the default constructor Test()
 cout << "t = " << t.getX() << endl;</pre>
 t.setX(4); calls the default Test(int x)
 cout << "t = " << t.getX() << endl;</pre>
 Test s(t); calls the copy constructor
  cout << "s = " << s.getX() << endl;</pre>
  return 0;
Answer: the output is
t = 0
t = 4
s = 5
```

c) [6 points] Explain the difference between pointers and references in C++. Provide an example if needed.

Answer: they are very similar when passed as arguments to the function. In both cases they pass the address of the variable, and allow changing its value.

Difference are:

- A pointer can be NULL, reference must always refer to some variable
- A pointer can be reassigned to point somewhere else, reference cannot
- Pointers are a variable themselves (tamking space in the memory), references are not

Problem 2 [25 points]

A **Doubly Linked List** is a linked list where each node has a pointer to the next element and to the previous element.

```
struct DLL node {
  int value;
  struct DLL node* next;
  struct DLL node* prev;
typedef struct DLL node DLL node t;
typedef struct {
  DLL node t* head; // pointer to the first node
  DLL node t* tail; // pointer to the last node
} DLL t;
Implement the following functions for Doubly Linked List of ints.
The running time in parts a-c of must be O(1).
 a) [6 points] Write a function that adds a new node with a given value to the head of the list.
 void add to head() {DLL t* list, int value) {
   DLL node* new node = (DLL node*) malloc(sizeof(DLL node));
   if (!new_node) return; // malloc fail
   new_node->prev = NULL;
   new node->next = list->head;
   new node->value = value;
   list->head = new node; // update the head pointer
   if (list->tail != NULL) // if list was not empty
    node node->next->prev = new node; // update the previous head
   else // if list was empty
    list->tail = new node;
 b) [6 points] Write a function that removes the first node of the list, and returns its value.
 void remove from head() {DLL t* list) {
   DLL node* prev head = list->head;
   list->head = prev head->next; // update head
   if (list->head == NULL) // if list becomes empty, update the tail
      list->tail = NULL;
      list->head->prev = NULL;
   free(prev_head);
```

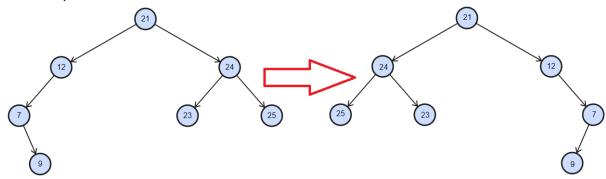
```
c) [6 points] Write a function that removes a given node from the list.
// assumption node indeed belongs to the list
int remove_node() {DLL_t* list, DLL_node_t* node) {
  // update node->prev -- the vertex before node
  if (node->prev != NULL) // node is not the first in the list
     node->prev->next = node->next;
  else // node is the first is in the list - update the head accordingly
     list->head= node->next:
  // update node->next -- the vertex after node
  if (node->next != NULL) // node is the last in the list
     node->next->prev = node->prev;
  else // node is last is in the list - update the tail accordingly
     list->tail = node->prev;
 int ret = node->value;
free(node);
 return ret;
d) [7 points] Write a function that gets a Doubly Linked List and a predicate p. The function
removes all nodes for which p(node->value)==false.
For example, if we apply it on the list head \rightarrow [1 \leftarrow -2 \leftarrow -5 \leftarrow -4 \leftarrow -8] \leftarrow tail with p=is_even(),
then the remaining list should be head \rightarrow [2 \leftarrow \rightarrow 4 \leftarrow \rightarrow 8] \leftarrow tail.
*Remember to release the memory of the deleted nodes.
void filter() {DLL t* list, bool(*p)(int)) {
  DLL node* cur = list->head;
  DLL node* next;
  while(cur) {
     next = cur->next; // save cur->next in case cur will be removed
     if (!p(cur->value))
        remove node(list, cur);
     cur = next; // move forward in the list
```

Problem 3 [25 points]

In this problem use the following struct representing a node in a Binary Tree of ints.

```
struct BTnode {
  int value;
  struct BTnode* left;
  struct BTnode* right;
  struct BTnode* parent;
};
typedef struct BTnode BTnode t;
```

a) [7 points] Write an algorithm that gets a **Binary Tree** and converts it into its mirror reverse. For example



```
void mirror tree(BTnode t* root) {
```

// idea: swap left and right child, and then apply recursion on each of the subtrees // it is possible to apply recursion first, and then swap the children

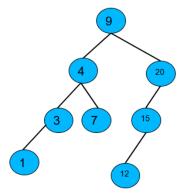
```
if (!root)
return;
```

```
// swap the two children (possibly one of them is NULL)
LL_node* tmp = root->left;
root->left = root->right;
root->right = tmp;

mirror_tree(root->left);
mirror_tree(root->right);
}
```

[3 points] Explain the running time of your function. Each node is processed exactly once for O(1) time. Therefore, the total time is O(size of tree)

b) [12 points] Write a function in C that gets a pointer to a node in a **Binary Search Tree**, and finds its predecessor. If the node is the minimal element in the tree, the function will return NULL. In the tree below: the predecessor of 4 is 3, the predecessor of 12 is 9, the predecessor of 7 is 4.



```
BTnode t* find predecessor(BTnode t* node) {
```

// idea: if node has a left child, then the predecessor is the largest node in the node->left. // otherwise, we can find the predecessor by going up "to the left".

```
if (node==NULL)
  return NULL;
```

```
if (node->left) { // node has a left child
    // find maximal node in the subtree node->left
BTnode_t* cur = node->left;
    while (cur->next != NULL) // iterate to the max in the subtree of node>tree
        cur = cur->right;
    return cur;
}
else { // node doesn't have a left child
    BTnode_t* cur = node;
    while (cur->parent && cur==cur>parent->left) // while cur is the left child of it parent
        cur = cur->parent;
    // here cur is either the right child of its parent or the root
    return cur->parent;
}
```

[3 points] Explain the running time of your function.

Answer: each of the loops goes either up the tree or down the tree once.

Therefore, the total running time is at most O(depth of tree)

Problem 4 [25 points]

```
a) [12 points] Write a function that gets an array of ints of length n, and outputs the length of
the longest contiguous increasing subsequence in O(n) time. For example, on input
[6,2,5,3,6,8,9,1] the answer should be 4, as the longest increasing subsequence is [3,6,8,9].
* If your solution runs in quadratic time or slower, you will get 8 points.
int longest incr subsequence(const int* arr, int n) {
 // idea: iterate through the array if arr[i]>arr[i-1], increase the count. Otherwise, reset count.
 // update max if needed
  if (n == 0)
    return 0;
  int i:
  int count = 1;
  int max = 1;
  for (i=1; i<n; i++) {
    // we allow equalities or strictly increasing - both versions accepted for full marks
    if( arr[i] >= arr[i-1] ) {
        count++;
        if (count>max) // can be made more efficient, but it doesn't affect O() complexity
           max = count;
     else // reset count to 1, because
        count = 1;
  return max;
[3 points] Explain the running time of your function.
We have a for loop with O(1) operations in each iteration. Therefore, the total time is O(n).
```

```
b) [10 points ] Write a function that gets an array of digits (given as ints) of length n, and
returns a string containing the largest number possible using these digits. For example:
       on input [6, 7, 5, 3, 8, 3, 0, 0] the function should return "87653300".
       on input [1, 2, 3, 4, 4] the function should return "44321".
       on input [0, 0, 0, 0, 0, 0] the function should return "0"
* You may assume all numbers in the array are between 0 and 9.
* Note: you need to return the number in a string because it may be too large to fit in an
int/lona.
** Remember to use dynamic memory allocation properly.
char* print max number(const int* arr, int n) {
// idea: count the number of 0's, 1's, 2's....and create a string using this statistics
  int count digits[10]; // count digits is created on the stack - no need to free it in the end
  int i,j,k;
  for (i=0; i<10; i++)
    count digits[i] = 0; // important to initialize;
 for (k=0; k<n; k++)
     count digits[arr[k]]++;
 // here count digits[i] contains the number of i's in the array;
  if (count_digits[0] == n) { // edge case of all zeros
    char* ret = (char*) malloc(2);
    ret[0] = '0'; ret[1] = '\0';
    return ret;
char* max_number = (char*) malloc(n+1); // +1 for string terminator '\0'
  // populate max_number - probably there is a cleaner way to write it
  k = 0;
  for (i=9; i>=0; i--)
    for (j=0; j<count_digits[i]; j++) {
      max number[k] = '0' + i;
      k++;
  max_number[n] = '\0';
  return max_number;
```