

# **Theorems and Proofs II**

## Previous Lecture

- *Axiomatic systems*
- *Informal proof*
- *Proof methods*
  - *Direct proof*
  - *Proof by contrapositive*

## Proof = logical inference

- Any theorem can be expressed in symbolic logic
- The formula being proved gives a (direct, mechanical) proof structure

Formula	Direct inference
$\forall x P(x)$	Show $P(a)$ for an arbitrary $a$
$\exists x P(x)$	Show $P(a)$ for some specific $a$
$P \rightarrow Q$	Assume $P$ and show $Q$
$P \leftrightarrow Q$	Show $P \rightarrow Q$ and $Q \rightarrow P$
$P \wedge Q$	Show $P$ and $Q$
$P \vee Q$	Show $P$ or show $Q$

- Can always rewrite to a **logically equivalent formula** and prove that  
e.g. Proof by contrapositive
- What about other logical inferences?

## Proof by Contradiction

- Based on the *rule of contradiction*

$$\frac{\neg p \rightarrow F}{\therefore p}$$

- Can be used to prove propositions of any form

- Proof structure:

1. Assume  $\neg p$ .
2. Infer a contradiction, i.e.  $F$
3. Conclude  $p$ .

*Usually by  
showing  $r$  and  $\neg r$*



## Example

- **Definition:** *a barber is called **strict** if he shaves those and only those who do not shave themselves.*
- **Theorem.** *There is no strict barber.*
- **Proof.**
- *Assume the contrary: a strict barber ( $c$  such that  $S(c)$ ) exists*
- *Does he shave himself – say  $Q(c)$ ?*
- *If no ( $\neg Q(c)$ ), then by definition he must shave himself ( $Q(c)$ )*
- *If yes ( $Q(c)$ ), then by definition he must not ( $\neg Q(c)$ )*
- *Either way we have  $Q(c) \wedge \neg Q(c)$ , a contradiction*
- *We conclude that a strict barber does not exist*

## Example (cntd)

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## Another Example

● *Definition:* a real number is said to be **rational** if it can be represented as a fraction  $\frac{a}{b}$  where  $a, b$  are integers

● *Prove that  $\sqrt{2}$  is irrational*

● *Proof*

*Suppose that  $\sqrt{2}$  is rational*

*Then there are integers  $a, b$  such that  $\sqrt{2} = \frac{a}{b}$ .*

*We may assume that  $a, b$  have no common divisor.*

*Squaring we obtain  $a^2 = 2b^2$ .*

*Since  $a^2$  is even,  $a$  is also even, hence  $a = 2c$  for some  $c$ .*

*Therefore  $2b^2 = 4c^2$ , and so  $b^2 = 2c^2$ .*

*Hence  $b$  is even.*

*We get that  $a$  and  $b$  have a common factor – 2. A contradiction.*

## Proof by cases

- Based on the inference rule

$$\frac{p \rightarrow r \quad q \rightarrow r}{\therefore (p \vee q) \rightarrow r}$$

- Used to prove implications where the premise *can be decomposed into a small number of cases*
- Proof structure (proving  $p \rightarrow q$ ):
  1. Assume  $p$
  2. *Decompose  $p$  into cases  $p_1 \vee p_2$  --- that is, infer  $p \rightarrow p_1 \vee p_2$*
  3. Show that  $p_1 \rightarrow q$  (Case 1)
  4. Show that  $p_2 \rightarrow q$  (Case 2)
  5. Conclude  $p \rightarrow q$ .

## Example

● *Theorem:*

*If  $n$  is an integer, then  $n^2 \geq n$*

● *Proof:*

*Suppose  $n$  is an integer*

*Then either (1)  $n = 0$ , (2)  $n > 0$ , or (3)  $n < 0$*

*Case 1:  $n = 0$*

*Then  $n^2 = 0 = n$ , so  $n^2 \geq n$*

*Case 2:  $n > 0$*

*Then  $n \geq 1$ , so multiplying both sides by  $n$ ,  $n^2 \geq n$*

*Case 3:  $n < 0$*

*Then  $n^2 \geq 0 > n$ , so  $n^2 \geq n$*

*In all cases,  $n^2 \geq n$*

*Q.E.D.*

## Proof play-by-play

- In the last example, the theorem had the form

$$\forall x (P(x) \rightarrow Q(x))$$

so why was this valid?

- We implicitly used that fact that

$$\forall x (\text{Integer}(x) \rightarrow (x = 0) \vee (x > 0) \vee (x < 0))$$

- Then we showed for arbitrary  $n$  **by case inference** that

$$(n = 0) \vee (n > 0) \vee (n < 0) \rightarrow n^2 \geq n$$

- Then we used **sylogism** to infer the conclusion:

$$\text{Integer}(n) \rightarrow (n = 0) \vee (n > 0) \vee (n < 0)$$

$$(n = 0) \vee (n > 0) \vee (n < 0) \rightarrow n^2 \geq n$$

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$$\therefore \text{Integer}(n) \rightarrow n^2 \geq n$$

## Exhaustive proof

- *Special case of proof by cases*
- *Replace an arbitrary  $a$  with disjunction of every possible value*

*Theorem:*

$$\text{if } 0 < n \leq 4, \text{ then } 3^n \leq (n + 1)^3$$

*Proof:*

*There are 4 cases to consider:  $n = 1, 2, 3,$  or  $4$*

*Case 1:  $n = 1$*

$$\text{Then } 3^1 = 3 \leq (1 + 1)^3 = 8$$

*Case 2:  $n = 2$*

$$\text{Then } 3^2 = 9 \leq (2 + 1)^3 = 27$$

*Case 3:  $n = 3$*

*“By calculation”*

*“remaining cases follow similarly”*

*Don't do this  
(proof by intimidation)*

*Q.E.D.*

## Excluding cases

- Common error is incorrectly thinking you considered all cases  
e.g. Case 1:  $x > 0$ , Case 2:  $x < 0$
- If two cases are *clearly identical*, can just show one

Theorem:

*if  $xy$  and  $x + y$  are even, then  $x$  and  $y$  are both even*

Proof:

Suppose  $x$  and  $y$  are not both even

*Without loss of generality*, assume  $x$  is odd

Case 1:  $y$  is even

Then  $x+y = 2k+1 + 2k' = 2(k+k') + 1$  is odd

Case 2:  $y$  is odd

Then  $xy = (2k+1)(2k'+1) = 2(2kk' + k + k') + 1$  is odd

*“other case follows identically”*

## Constructive vs non-constructive proofs

- *Want to prove the existential statement  $\exists x P(x)$ .*
- *Direct method involves **finding a witness  $c$**  such that  $P(c)$  is true.*  
*e.g. Prove that there exists a red car...*  
*My car is red!*
- *Non-constructive proofs do not find an explicit  $c$ . How?*
- *One method: proof by contradiction*

$$\begin{array}{ccc}
 \frac{\neg \exists x P(x) \rightarrow F}{\therefore \exists x P(x)} & \xrightarrow{\text{green arrow}} & \frac{\forall x \neg P(x) \rightarrow F}{\therefore \exists x P(x)}
 \end{array}$$

## Example

● *Theorem:*

*There exist irrational numbers  $x, y$  such that  $x^y$  is rational*

● *Proof:*

*Suppose instead that for all irrational  $x, y$ ,  $x^y$  is irrational*

*Consider the number  $\sqrt{2^{\sqrt{2}}}$ . Either it is rational or irrational*

*Case 1:  $\sqrt{2^{\sqrt{2}}}$  is rational*

*Then since  $x=y=\sqrt{2}$  is irrational we have a contradiction*

*Case 2:  $\sqrt{2^{\sqrt{2}}}$  is irrational*

*Then  $x=\sqrt{2^{\sqrt{2}}}$  and  $y = \sqrt{2}$  are both irrational*

*But  $x^y = (\sqrt{2^{\sqrt{2}}})^{\sqrt{2}} = \sqrt{2^{\sqrt{2}\sqrt{2}}} = \sqrt{2^2} = 2$  is rational*

*Hence we have a contradiction.*

*Q.E.D.*

# Summary

- *Proofs are logical arguments*
- *A valid proof must use valid logical inferences*
- *Knowing which statements are logically equivalent and valid rules of inferences gives access to a wide array of valid proofs*
- *If in doubt, write it out! (in symbolic, explicit form)*

# Practice

*Exercises from the Book:*

*7<sup>th</sup> edition: 3, 9, 15, 17 (page 91)*

*8<sup>th</sup> edition: 3, 9, 17, 19 (page 95)*