

Relations

Previous Lecture

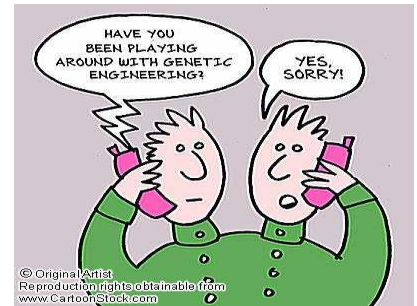
- *Set-based models of logic*
 - *Power sets & propositional logic*
 - *Binary representation of sets*
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- *Fun fact:*
 - there are 2^{2^n} unique truth tables in n propositional variables*
 - *Fun questions:*
 - What does implication correspond to?*
 - What are the models of a quantified statement?*

Relations

- The predicate $P(x,y) = x > y$ isn't a set of values in the same sense... it's a **relationship** that holds between the values of x and y

- Between people, familial relationships

`to be brothers`	x is a brother of y
`to be older`	x is older than y
`to be parents`	x and y are parents of z



- Between things, numerical relationships

`to be greater than`	$x < y$ on the set of real numbers
`to be divisible by`	x is divisible by y on the set of integers
- Between things and people, legal relationships

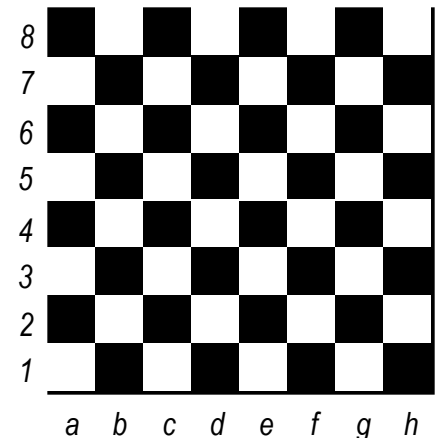
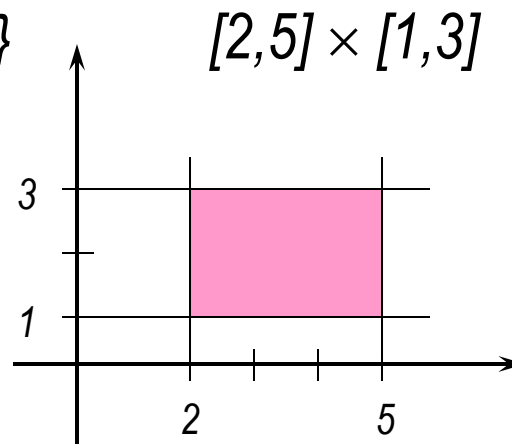
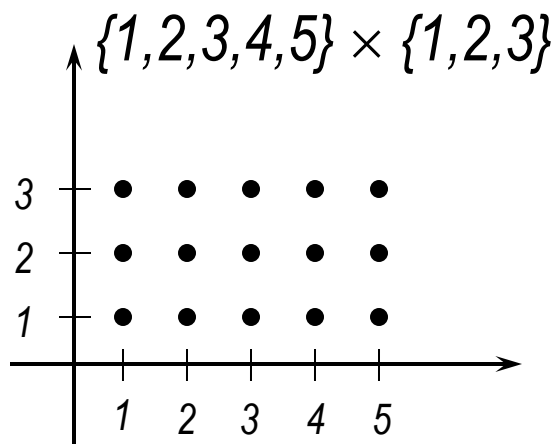
`to be an owner`	x is an owner of y
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Relations as sets: the Cartesian Product

- The **Cartesian product** of sets A and B , denoted by $A \times B$, is the set of all **ordered pairs** (a,b) of elements from A and B .

$$A \times B = \{ (a,b) \mid a \in A, b \in B \}$$

- The elements of the Cartesian product are ordered pairs. In particular, $(a,b) = (c,d)$ if and only if $a = c$ and $b = d$.
- If sets are thought of as '1-dimensional' objects, Cartesian products are 2-dimensional



Example

● *Are the following sets equal or not equal?*

$\{1,2,3\}$ and $\{3,2,1\}$?

$\{\{1,2,3\},\{3,2,1\}\}$ and $\{\{1,2,3\}\}$?

$\{(1,2,3)\}$ and $\{(3,2,1)\}$?

$\{(1,2,3), (3,2,1)\}$ and $\{(1,2,3)\}$?

Cardinality of Cartesian Product

● Theorem.

$$|A \times B| = |A| \cdot |B|$$

● *Proof*

Consider an element a of A

We can construct $|B|$ unique pairs (a,b) corresponding to the $|B|$ elements of B

Since we have $|A|$ such unique elements a , in total there are $|A| \cdot |B|$ unique ordered pairs.

Q.E.D.

Back to predicates and their meaning

- Consider the predicates

$$P(x,y) = x > y,$$

$$Q(x,y) = y > x$$

- They have 2 indeterminates, so they aren't quite sets of values in the same sense.
- They give a *relationship* between two indeterminates x and y , not just a set of values.
- Their models are hence *sets of ordered pairs*, i.e. a subset of the Cartesian product of the variable's domains called a *relation*

$$P(x,y) \equiv \{ (x,y) \mid P(x,y) \} \subseteq U \times U$$

Binary Relations

- A **binary relation** from set A to set B is any subset of $A \times B$.
If $A = B$ then we say that the relation is **on** the set A

' x is a brother of y ' \subseteq $People \times People$

' x is older than y ' \subseteq $People \times People$

' x is an owner of y ' \subseteq $People \times Properties$

' $x < y$ ' \subseteq $\mathbb{R} \times \mathbb{R}$

' x divides y ' \subseteq $\mathbb{Z} \times \mathbb{Z}$



- **Binary predicates** $P(x,y)$ correspond to binary relations!

$$P(x,y) \equiv \{ (x,y) \mid P(x,y) \} \subseteq U \times U$$

Sets, Relations, and Predicates

- *Now binary predicates have a direct correspondence to sets too!*

Unary predicate

$$P(x)$$

Set

$$A = \{ x \mid P(x) \}$$

Binary predicate

$$P(x,y)$$

Binary relation

$$R = \{ (x,y) \mid P(x,y) \}$$

Ternary predicate

$$P(x,y,z)$$

???

Cartesian Product of More Than Two Sets

- Instead of ordered pairs we may consider ordered **triples**, or, more general, **k-tuples**.

(a,b,c) , an ordered triple

(a,b,c,d) , an ordered quadruple

(a_1, a_2, \dots, a_k) a k-tuple

- Triples, quadruples, and k-tuples are elements of Cartesian products of 3, 4, and k sets, respectively

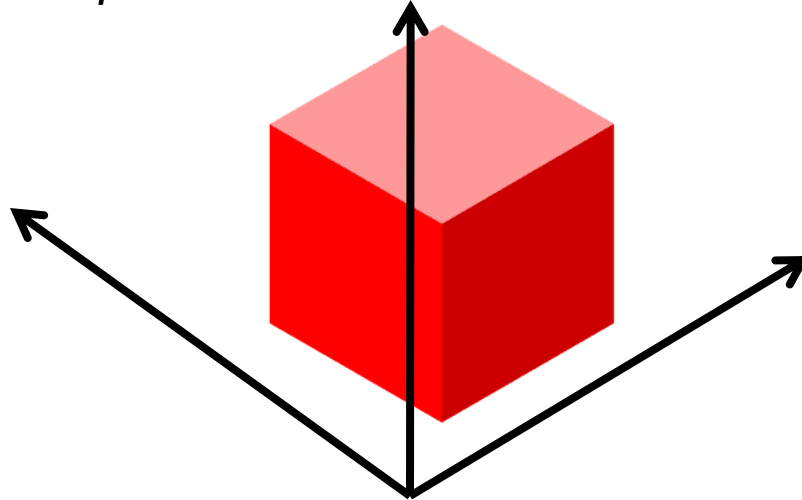
$$A \times B \times C = \{(a,b,c) \mid a \in A, b \in B, c \in C\}$$

$$A \times B \times C \times D = \{(a,b,c,d) \mid a \in A, b \in B, c \in C, d \in D\}$$

$$A_1 \times A_2 \times \dots \times A_k = \{(a_1, a_2, \dots, a_k) \mid a_1 \in A_1, a_2 \in A_2, \dots, a_k \in A_k\}$$

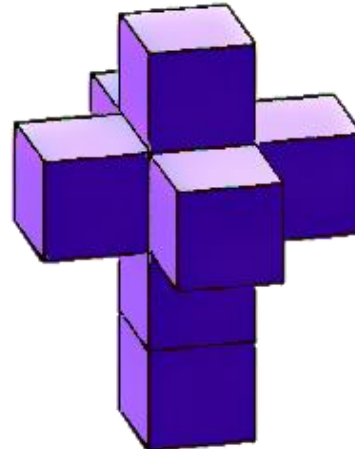
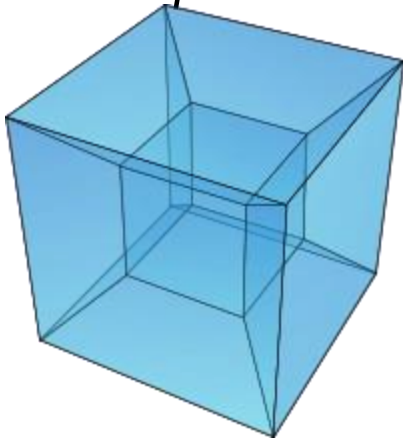
Cartesian Product of More Than Two Sets

- *Cartesian product of 3 sets can be viewed as a '3-dimensional' set*



*E.g. Coordinates
(a,b,c) in $R \times R \times R$*

- *Cartesian product of more than 3 sets is a multi-dimensional set*



More Relations (cntd)

● *Binary relations can be generalized to subsets of Cartesian products of more than two sets.*

● *A subset of the Cartesian product of 3 sets is a **ternary relation***

'x and y are parents of z' \subseteq People \times People \times People

● *Any subset of the Cartesian product of k sets is a **k-ary relation***

$\{(a_1, a_2, \dots, a_k) \mid a_1 + a_2 + \dots + a_k = 3\} \subseteq \mathbb{R} \times \mathbb{R} \times \dots \times \mathbb{R}$

Relational Databases

- A *relational database* is a collection of *tables* like

No.	Name	Student ID	Supervisor	Thesis title
1.	Bradley Coleman	30101234	Petra Berenbrink	Algebraic graph theory
...

A table consists of a schema and an instance ...

The instance of this table is a 5-ary relation, a subset of the Cartesian product

$$\mathbb{Z}^+ \times \text{Names} \times \text{8-strings_of_digits} \times \text{Names} \times \text{Meaningful_Sentences}$$

Cartesian Product, Intersection and Union

● **Theorem.** For any sets A, B, C

$$(1) A \times (B \cap C) = (A \times B) \cap (A \times C)$$

$$(2) A \times (B \cup C) = (A \times B) \cup (A \times C)$$

$$(3) (A \cap B) \times C = (A \times C) \cap (B \times C)$$

$$(4) (A \cup B) \times C = (A \times C) \cup (B \times C)$$

● *Proof (of (2))*

$$\begin{aligned} A \times (B \cup C) &= \{ (a,b) \mid a \in A \wedge b \in B \cup C \} \\ &= \{ (a,b) \mid a \in A \wedge (b \in B \vee b \in C) \} \\ &= \{ (a,b) \mid (a \in A \wedge b \in B) \vee (a \in A \wedge b \in C) \} \\ &= \{ (a,b) \mid a \in A \wedge b \in B \} \cup \{ (a,b) \mid a \in A \wedge b \in C \} \\ &= (A \times B) \cup (A \times C) \end{aligned}$$

Q.E.D.

Sets, Relations, and Predicates

- *So all of predicate logic can be modelled in Set theory*

Unary predicate

$$P(x)$$

Set

$$A = \{ x \mid P(x) \}$$

Binary predicate

$$P(x,y)$$

Binary relation

$$R = \{ (x,y) \mid P(x,y) \}$$

Ternary predicate

$$P(x,y,z)$$

Ternary relation

$$R = \{ (x,y,z) \mid P(x,y,z) \}$$

Is this a good idea?

Everything everywhere all at once

- Recall Russell's paradox U – the set of all sets which do not contain themselves
- Paradoxes are inconsistencies which lead to *proofs of untrue facts*
- Theorem: Let P be a logical formula. Then P is true
- Proof:

Let $U = \{ S \mid S \notin S \text{ i.e. } S \text{ is a set which does not contain itself} \}$

Case 1: $U \in U$

Then $U \notin U$ by definition, hence F is true

Case 2: $U \notin U$

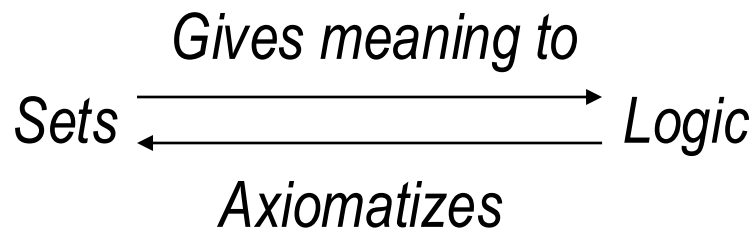
Then $U \in U$ by definition, hence F is true

Since F is true, P by the law of contradiction.

Q.E.D.

A history lesson

- *Frege was writing his formalization of mathematics based on logic when Russell sent him the paradox in 1901*
- *The paradox worked because of assumed **extensionality** --- the existence of set-theoretic models which characterize predicates*
- *One method of fixing was by consistently axiomatizing set theory **in logic** (ZFC Set Theory), using universes (championed by **Von Neumann**) to avoid paradoxes*
- *The result is **circular** or **impredicative**:*



Intuitionism

- Competing “fix” via Brouwer was to *reject logic altogether*
- At the time, logic was a tool for discovering the objective truth of a mathematical statement
- Intuitionists were Nihilists who said that
“there is no objective truth, only a mathematician’s intuition”
- Meanwhile Russell disagreed with everyone:
“Logic exists, but should not be circular”
- Heyting (1930’s) committed blasphemy (at the time) and developed *intuitionistic logic*:
“The logic of what can be proved, not objective truth”

Intuitionistic logic

- *Formulas have no intrinsic truth*
- *P means “ P is provable”*
- *$\neg P$ is “unprovability” --- that is, provability of P implies a contradiction*
- *Law of the excluded middle is not valid:
 $P \vee \neg P$ is “ P is provable, or the unprovability of P is provable”*
- *Double negation is also not valid:
 $\neg \neg P$ means ???*
- *BHK interpretation:
“Meaning” of a formula is **A COMPUTER PROGRAM***

Practice

Exercises from the Book:

7th edition: 27, 32, 38, 39 (page 126)

8th edition: 29, 34, 40, 41 (page 132 – 133)