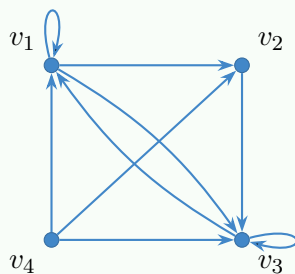


Problem 1. Draw a graph with the adjacency matrix

$$\begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 \end{pmatrix}.$$

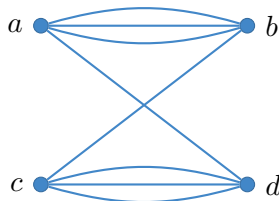
Answer.

We use the convention that the entry in row i and column j is the number of directed edges from v_i to v_j . One possible drawing is



For example, the first row gives the loop at v_1 and the edges $v_1 \rightarrow v_2$ and $v_1 \rightarrow v_3$. The zero fourth column means that no edge enters v_4 .

Problem 2. Represent the following graph using adjacency lists, an adjacency matrix, and an incidence matrix.



Answer.

Repeated entries in an adjacency list record parallel edges. Thus

$$\begin{aligned} a &: (b, b, b, d), & b &: (a, a, a, c), \\ c &: (d, d, d, b), & d &: (c, c, c, a). \end{aligned}$$

With the vertex order a, b, c, d , the adjacency matrix is

$$A = \begin{pmatrix} 0 & 3 & 0 & 1 \\ 3 & 0 & 1 & 0 \\ 0 & 1 & 0 & 3 \\ 1 & 0 & 3 & 0 \end{pmatrix}.$$

Let e_1, e_2, e_3 be the three ab -edges, let e_4, e_5, e_6 be the three cd -edges, and let $e_7 = ad$ and

$e_8 = bc$. The incidence matrix is

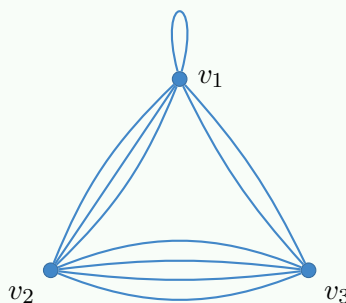
	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
a	1	1	1	0	0	0	1	0
b	1	1	1	0	0	0	0	1
c	0	0	0	1	1	1	0	1
d	0	0	0	1	1	1	1	0

Problem 3. Draw an undirected multigraph represented by the adjacency matrix

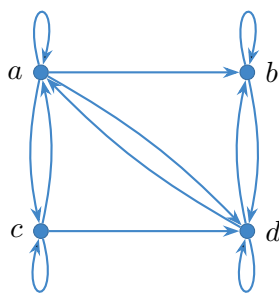
$$\begin{pmatrix} 1 & 3 & 2 \\ 3 & 0 & 4 \\ 2 & 4 & 0 \end{pmatrix}.$$

Answer.

The diagonal entry 1 gives one loop at v_1 . The other entries give three v_1v_2 -edges, two v_1v_3 -edges, and four v_2v_3 -edges.



Problem 4. Represent the following directed multigraph using adjacency lists, an adjacency matrix, and an incidence matrix.



Answer.

We list outgoing neighbours. A loop therefore places the vertex itself in its list:

$$\begin{aligned} a &: (a, b, c, d), & b &: (b, d), \\ c &: (a, c, d), & d &: (a, b, d). \end{aligned}$$

With vertex order a, b, c, d , the adjacency matrix is

$$A = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \end{pmatrix}.$$

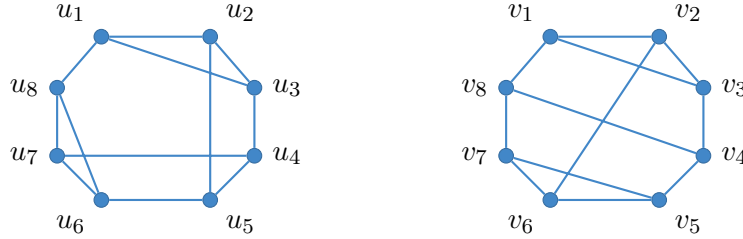
For the incidence matrix, use -1 at the tail and $+1$ at the head. Let e_1, e_2, e_3, e_4 be the loops at a, b, c, d , respectively, and set

$$\begin{aligned} e_5 &= a \rightarrow b, & e_6 &= c \rightarrow d, & e_7 &= a \rightarrow c, & e_8 &= c \rightarrow a, \\ e_9 &= b \rightarrow d, & e_{10} &= d \rightarrow b, & e_{11} &= a \rightarrow d, & e_{12} &= d \rightarrow a. \end{aligned}$$

A loop has the same head and tail, so in the incidence column we put a ± 1 . Thus

	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8	e_9	e_{10}	e_{11}	e_{12}
a	± 1	0	0	0	-1	0	-1	1	0	0	-1	1
b	0	± 1	0	0	1	0	0	0	-1	1	0	0
c	0	0	± 1	0	0	-1	1	-1	0	0	0	0
d	0	0	0	± 1	0	1	0	0	1	-1	1	-1

Problem 5. Use graph parameters to determine whether the following graphs are isomorphic.



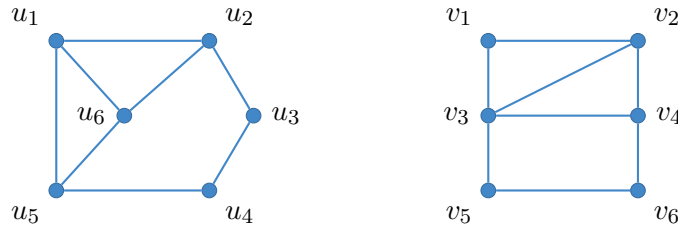
Answer.

Both graphs have 8 vertices, 12 edges, and degree sequence $(3, 3, 3, 3, 3, 3, 3, 3)$. They are in fact isomorphic. One isomorphism is

$$\begin{array}{c|cccccccc} u_i & u_1 & u_2 & u_3 & u_4 & u_5 & u_6 & u_7 & u_8 \\ \hline \varphi(u_i) & v_2 & v_1 & v_3 & v_4 & v_8 & v_7 & v_5 & v_6 \end{array}.$$

Checking the outer-cycle edges and the four chords shows that every edge is sent to an edge, so φ is an isomorphism.

Problem 6. Use graph parameters to determine whether the following graphs are isomorphic.



Answer.

The degree sequences are different:

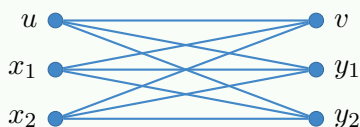
$$(3, 3, 3, 3, 2, 2) \quad \text{and} \quad (4, 3, 3, 2, 2, 2).$$

In particular, the second graph has a vertex of degree 4, while the first does not. Since an isomorphism preserves degrees, the graphs are not isomorphic.

Problem 7. Find the number of paths of lengths 2, 3, 4, and 5 between any two adjacent vertices in $K_{3,3}$.

Answer.

The graph $K_{3,3}$ is the following (the label of vertices doesn't matter too much).



Let the two parts be

$$L = \{u, x_1, x_2\}, \quad R = \{v, y_1, y_2\}.$$

Even-length paths cannot go from $u \in L$ to $v \in R$, because an even path starts and ends at the same side of the graph.

A path of length 3 has the form

$$u \longrightarrow y_i \longrightarrow x_j \longrightarrow v,$$

giving $2 \cdot 2 = 4$ choices.

A path of length 5 has the form

$$u \longrightarrow y_i \longrightarrow x_j \longrightarrow y_k \longrightarrow x_\ell \longrightarrow v.$$

There are two choices each for y_i and x_j , after which y_k and x_ℓ are forced. Thus, there are also $2 \cdot 2 = 4$ choices for paths of length 5 as well.

Problem 8. Show that in every simple graph there is a path from every vertex of odd degree to some other vertex of odd degree.

Answer.

Let v have odd degree, and consider the connected component C containing v . By the handshaking lemma, C has an even number of odd-degree vertices. Since v is one of them, C contains at least one other odd-degree vertex w . The vertices v and w lie in the same connected component, so there is a path from v to w .

Problem 9. How many nonisomorphic simple graphs with 4 vertices are there?

Answer.

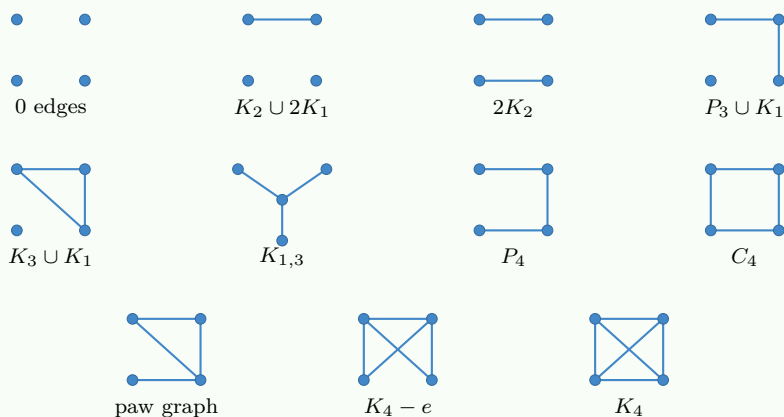
Classify the graphs by their number m of edges. Taking complements shows that the counts for m and $6 - m$ are equal. The counts are

m	0	1	2	3	4	5	6
number of types	1	1	2	3	2	1	1

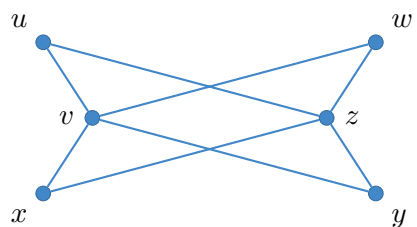
and hence the total is

$$1 + 1 + 2 + 3 + 2 + 1 + 1 = 11.$$

The following are representatives of the eleven types.



Problem 10. Is the following graph bipartite?



Answer.

Yes. A bipartition is

$$\{v, z\} \quad \text{and} \quad \{u, x, w, y\}.$$

Each of u, x, w, y is adjacent to both v and z , and there are no other edges. Thus the graph is in fact $K_{2,4}$.